

Information-theoretic thresholds

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based on joint work with

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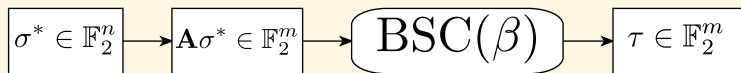
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Inference from samples

- ▶ to infer an unknown probability distribution from samples
- ▶ the distribution itself is random, determined by parameters σ^*

Example: error-correcting codes



- ▶ $A \in \mathbb{F}_2^{m \times n}$ is the generator matrix
- ▶ $A\sigma^*$ is subjected to noise

Example: the stochastic block model

	$\sigma^{*-1}(1)$	$\sigma^{*-1}(2)$	$\sigma^{*-1}(3)$
$\sigma^{*-1}(1)$	$e^{-\beta}$	1	1
$\sigma^{*-1}(2)$	1	$e^{-\beta}$	1
$\sigma^{*-1}(3)$	1	1	$e^{-\beta}$

- ▶ random coloring $\sigma^* : V \rightarrow \{1, \dots, q\}$
- ▶ for each $e = \{v, w\}$ independently,

$$\mathbb{P}[e \in \mathbf{G}^* | \sigma^*] = \frac{d}{n} \cdot \frac{q}{q-1+e^{-\beta}} \cdot \begin{cases} e^{-\beta} & \text{if } \sigma^*(v) = \sigma^*(w), \\ 1 & \text{if } \sigma^*(v) \neq \sigma^*(w) \end{cases}$$

- ▶ d = signal strength; $e^{-\beta}$ = noise

Example: the stochastic block model

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- ▶ the agreement of $\sigma, \tau : V \rightarrow \{1, \dots, q\}$ is

$$\alpha(\sigma, \tau) = \frac{1}{q-1} \max_{\kappa \in S_q} \left\{ \frac{q}{n} \sum_{v \in V} \mathbf{1}\{\sigma(v) = \kappa \circ \tau(v)\} - 1 \right\}.$$

- ▶ for what d, β is it possible to recover $\tau_{\mathbf{G}^*}$ such that

$$\mathbb{E}[\alpha(\sigma^*, \tau_{\mathbf{G}^*})] \geq \Omega(1) ?$$

Example: the stochastic block model

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Easy-hard-impossible

- ▶ for large d efficient algorithms should detect σ^*
- ▶ for very small d there is nothing to detect
- ▶ in-between the problem may be well-posed but hard

$$0 < d_{\text{inf}}(\beta) < d_{\text{alg}}(\beta)$$

Example: the stochastic block model

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The algorithmic threshold

- ▶ combinatorial algorithms for large d [1980s]
- ▶ spectral algorithms for moderate d [1990s, 2000s]
- ▶ the Kesten-Stigum threshold [AS15]

$$d_{\text{alg}}(\beta) \stackrel{?}{=} \left(\frac{q-1+e^{-\beta}}{1-e^{-\beta}} \right)^2$$

Example: the stochastic block model

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The information-theoretic threshold

- ▶ statistical physics prediction
- ▶ the case $q = 2$
- ▶ bounds on $d_{\text{inf}}(q, \beta)$

[DKMZ11]

[MNS13, MNS14, M14]

[BMNN16]

The information-theoretic threshold

Theorem

[COKPZ16]

For $\beta > 0$, $d > 0$ let

$$\mathcal{B}_{q,\beta}^*(d) = \sup \left\{ \mathcal{B}_{q,\beta,d}(\pi) : \pi = \mathcal{T}_{q,\beta,d}(\pi), \int \mu(i) d\pi(\mu) = 1/q \right\} \quad \text{where}$$

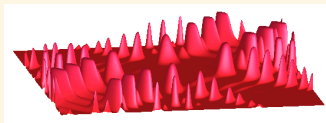
$$\mathcal{T}_{q,\beta,d} : \pi \mapsto \sum_{\gamma=0}^{\infty} \frac{d^\gamma \exp(-d)}{q(1 - (1 - e^{-\beta})/q)^\gamma \gamma!} \int \left[\sum_{h=1}^q \prod_{j=1}^{\gamma} 1 - (1 - e^{-\beta})\mu_j(h) \right] \delta_{\text{BP}_{\mu_1, \dots, \mu_\gamma}} d\pi^{\otimes \gamma}(\mu_1, \dots, \mu_\gamma),$$

$$\text{BP}_{\mu_1, \dots, \mu_\gamma}(i) = \frac{\prod_{j=1}^{\gamma} 1 - (1 - e^{-\beta})\mu_j(i)}{\sum_{h=1}^q \prod_{j=1}^{\gamma} 1 - (1 - e^{-\beta})\mu_j(h)},$$

$$\mathcal{B}_{q,\beta,d}(\pi) = \mathbb{E} \left[\frac{\Lambda(\sum_{\sigma=1}^q \prod_{i=1}^{\gamma} 1 - (1 - e^{-\beta})\mu_i^{(\pi)}(\sigma))}{q(1 - (1 - e^{-\beta})/q)^\gamma} - \frac{d}{2} \frac{\Lambda(1 - (1 - e^{-\beta}) \sum_{\sigma=1}^q \mu_1^{(\pi)}(\sigma) \mu_2^{(\pi)}(\sigma))}{1 - (1 - e^{-\beta})/q} \right].$$

Then $d_{\text{inf}}(q, \beta) = \inf \left\{ d > 0 : \mathcal{B}_{q,\beta}^*(d) > \ln q + \frac{d}{2} \ln(1 - (1 - e^{-\beta})/q) \right\}$.

The posterior distribution



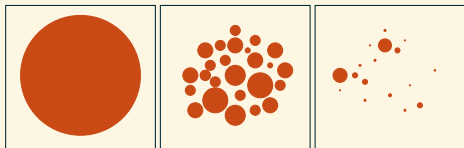
- ▶ define

$$\psi_{\mathbf{G}^*}(\sigma) = \prod_{\{v,w\} \in E(G)} \exp(-\beta \mathbf{1}\{\sigma(v) = \sigma(w)\}),$$

$$Z(\mathbf{G}^*) = \sum_{\sigma \in \Omega^V} \psi_{\mathbf{G}^*}(\sigma).$$

- ▶ then $\mathbb{P}[\boldsymbol{\sigma}^* = \sigma | \mathbf{G}^*] \asymp \mu_{\mathbf{G}^*}(\sigma) = \psi_{\mathbf{G}^*}(\sigma) / Z(\mathbf{G}^*)$

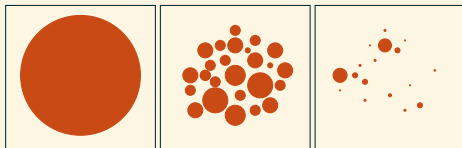
The posterior distribution



- ▶ reconstruction is impossible iff

$$\lim_{n \rightarrow \infty} \frac{1}{n^2} \sum_{v,w} \mathbb{E} \|\mu_{\mathbf{G}^*,v,w} - \mu_{\mathbf{G}^*,v} \otimes \mu_{\mathbf{G}^*,w}\|_{\text{TV}} = 0$$

The posterior distribution

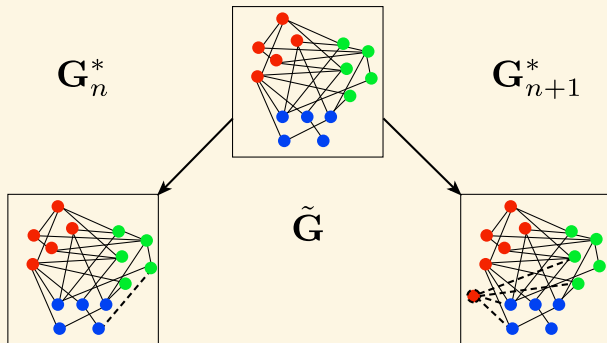


$$\lim_{n \rightarrow \infty} \frac{1}{n^2} \sum_{v,w} \mathbb{E} \|\mu_{\mathbf{G}^*,v,w} - \mu_{\mathbf{G}^*,v} \otimes \mu_{\mathbf{G}^*,w}\|_{\text{TV}} = 0$$
$$\Leftrightarrow \lim_{n \rightarrow \infty} \frac{1}{n} \mathbb{E}[\log Z(\mathbf{G}^*)] = \log q + \frac{d}{2} \log(1 - (1 - e^{-\beta})/q).$$

The Aizenman-Sims-Starr scheme

$$\lim_{n \rightarrow \infty} \frac{1}{n} \mathbb{E}[\log Z(\mathbf{G}^*)] = \lim_{n \rightarrow \infty} \mathbb{E} \left[\log \frac{Z(\mathbf{G}_{n+1}^*)}{Z(\mathbf{G}_n^*)} \right]$$

The Aizenman-Sims-Starr scheme



$$\frac{Z(\tilde{G} + \overline{vw})}{Z(\tilde{G})} = \sum_{\sigma, \tau \in [q]} e^{-\beta \mathbf{1}\{\sigma \neq \tau\}} \mu_{G^*, v, w}(\sigma, \tau)$$

Correlations

- ▶ \mathcal{X} = fixed finite set
- ▶ $\mu \in \mathcal{P}(\mathcal{X}^n)$ for some large integer n

Correlations

Definition

A probability measure $\mu \in \mathcal{P}(\mathcal{X}^n)$ is ε -symmetric if

$$\frac{1}{n^2} \sum_{i,j=1}^n \|\mu_{i,j} - \mu_i \otimes \mu_j\|_{\text{TV}} < \varepsilon$$

Correlations

The magic lemma

[COKPZ16]

For any $\varepsilon > 0$ there is a **bounded** random variable T such that for all $\mu \in \mathcal{P}(\mathcal{X}^n)$ the following is true:

- ▶ choose $\mathbf{U} \subset \{1, \dots, n\}$ of size T randomly
- ▶ sample $\hat{\sigma}$ from μ
- ▶ let

$$\hat{\mu}(\tau) = \mu[\tau | \forall i \in \mathbf{U} : \tau(i) = \hat{\sigma}(i)];$$

then

$$\mathbb{P}[\hat{\mu} \text{ is } \varepsilon\text{-symmetric}] > 1 - \varepsilon$$

Correlations

Lemma

[BCO15]

For any $\varepsilon > 0$, $k \geq 3$ there is $\delta > 0$ s.t. for $n > 1/\delta$ for δ -symmetric μ ,

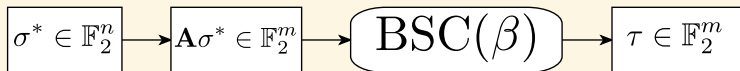
$$\frac{1}{n^k} \sum_{i_1, \dots, i_k=1}^n \|\mu_{i_1, \dots, i_k} - \mu_{i_1} \otimes \dots \otimes \mu_{i_k}\|_{\text{TV}} < \varepsilon$$

Low density generator matrix codes



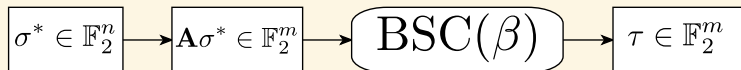
- ▶ $\mathbf{A} \in \mathbb{F}_2^{m \times n}$ with $k \geq 3$ ones per row
- ▶ signal $d = km/n$, noise β

Low density generator matrix codes



$$I(\sigma^*, \tau | A) = \sum_{s,t} P[\sigma^* = s, \tau = t] \log \frac{P[\sigma^* = s, \tau = t]}{P[\sigma^* = s] P[\tau = t]}$$

Low density generator matrix codes



- ▶ non-rigorous statistical physics analysis [KS99]
- ▶ upper bound on the mutual information, even k [M05]
- ▶ existence of $\lim_{n \rightarrow \infty} \frac{1}{n} I(\sigma^*, \tau | \mathbf{A})$, even k [AM15]

Low density generator matrix codes

Theorem

[CKPZ16]

For $k \geq 2$, $\beta > 0$, $d > 0$ and $\pi \in \mathcal{P}_0([-1, 1])$ let

$$\mathcal{B}_{d,\beta}(\pi) = \mathbb{E} \left[\frac{1}{2} \Lambda \left(\sum_{\sigma=\pm 1} \prod_{i=1}^r 1 + (1-2\beta)\sigma \mathbf{J}_i \prod_{j=1}^{k-1} \boldsymbol{\theta}_{i,j}^\pi \right) - \frac{d(k-1)}{k} \Lambda \left(1 + (1-2\beta) \mathbf{J} \prod_{j=1}^k \boldsymbol{\theta}_j^\pi \right) \right]$$

Then

$$\lim_{n \rightarrow \infty} \frac{1}{n} I(\boldsymbol{\sigma}^*, \boldsymbol{\tau} | \mathbf{A}) = (1 + d/k) \log 2 + \beta \log \beta + (1 - \beta) \log(1 - \beta) - \sup_{\pi \in \mathcal{P}_0([-1, 1])} \mathcal{B}_{d,\beta}(\pi)$$

The information-theoretic threshold is equal to

$$d_{\text{inf}}(\beta) = \inf \left\{ d > 0 : \sup_{\pi \in \mathcal{P}_0([-1, 1])} \mathcal{B}_{d,\beta}(\pi) > \log 2 \right\}$$

Conclusions

- ▶ generalisation: the “teacher-student scheme”
- ▶ justification of the ‘replica symmetric cavity method’
- ▶ other applications:
 - ▶ random graph colouring
 - ▶ Goldreich’s one-way function
 - ▶ the diluted p -spin model