

Which alternating dimaps are binary functions?

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23 June 2014

Cutset space

Incidence matrix of graph G :

$$\begin{array}{c} \text{vertices} \end{array} \begin{array}{c} \left(\begin{array}{c} \vdots \\ \dots \quad 0/1 \text{ entries} \quad \dots \\ \vdots \end{array} \right) \end{array} \begin{array}{c} \text{edges} \end{array}$$

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Cutset space := rowspace of incidence matrix over $GF(2)$.

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$f : 2^E \rightarrow \{0, 1\}$, defined by:

$$f(X) = \begin{cases} 1, & \text{if } X \text{ is in cutset space;} \\ 0, & \text{otherwise.} \end{cases}$$

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Often think of this as a vector, \mathbf{f} , length $2^{|E|}$, entries indexed by subsets of E (or their characteristic vectors).

Binary functions

Indicator functions of cutset spaces are prototypical *binary functions*.

Let E be a finite set (the *ground set*).

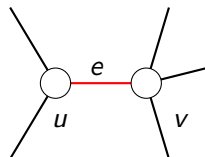
A **binary function** is a function $f : 2^E \rightarrow \mathbb{C}$ such that $f(\emptyset) = 1$.

In terms of vectors: it's a $2^{|E|}$ -element column vector \mathbf{f} , with entries indexed by subsets of E (or their characteristic vectors), such that $f_\emptyset = 1$.

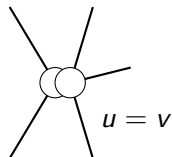
Back to graphs ...

Contraction and Deletion

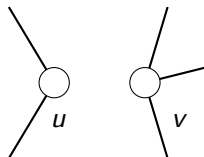
G



G/e



$G \setminus e$



Minors

H is a **minor** of G if it can be obtained from G by some sequence of deletions and/or contractions.

The order doesn't matter. Deletion and contraction **commute**:

$$\begin{aligned}G/e/f &= G/f/e \\ G \setminus e \setminus f &= G \setminus f \setminus e \\ G/e \setminus f &= G \setminus f/e\end{aligned}$$

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Importance of minors:

- ▶ excluded minor characterisations
 - ▶ planar graphs (Kuratowski, 1930; Wagner, 1937)
 - ▶ graphs, among matroids (Tutte, PhD thesis, 1948)
 - ▶ Robertson-Seymour Theorem (1985–2004)
- ▶ counting
 - ▶ Tutte-Whitney polynomial family

Duality and minors

Classical duality for embedded graphs:

$$\begin{array}{ccc} G & \longleftrightarrow & G^* \\ \text{vertices} & \longleftrightarrow & \text{faces} \end{array}$$

Duality and minors

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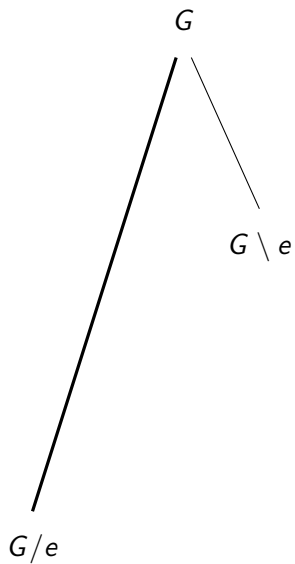
$$G \longleftrightarrow G^*$$

vertices \longleftrightarrow faces

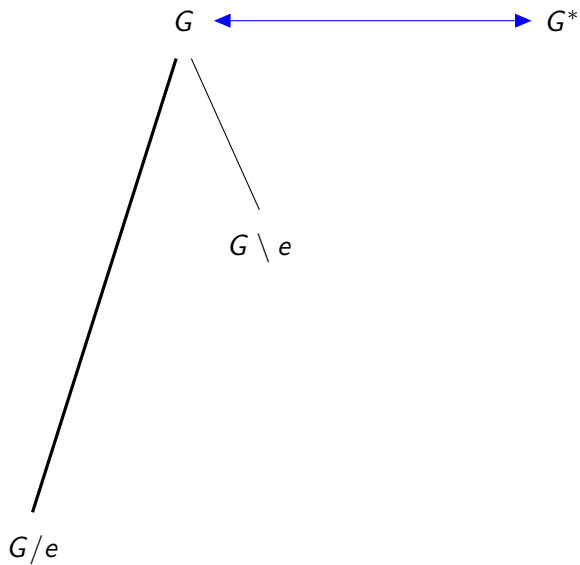
contraction \longleftrightarrow deletion

$$(G/e)^* = G^* \setminus e$$
$$(G \setminus e)^* = G^*/e$$

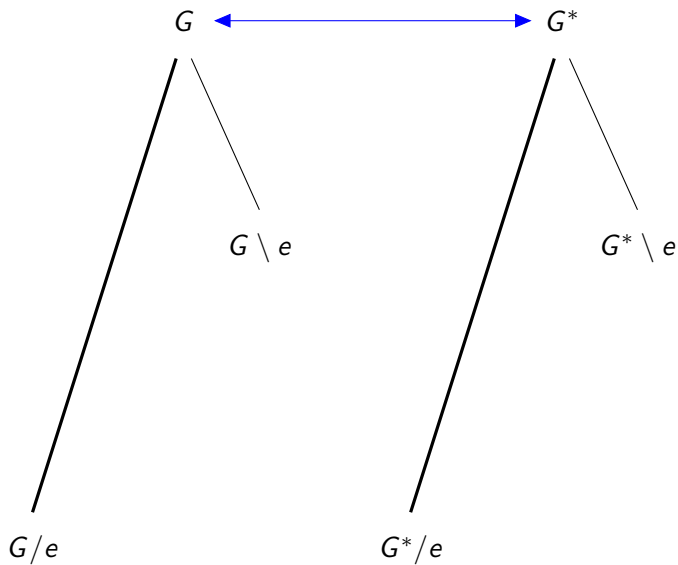
Duality and minors



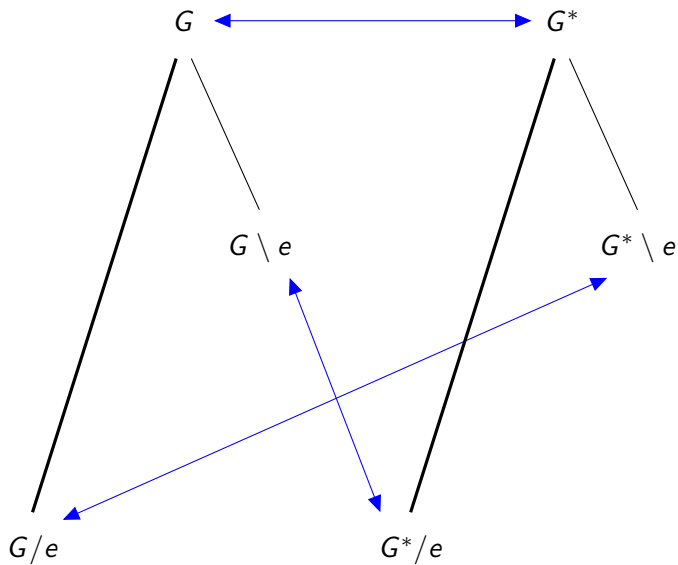
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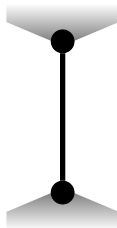


Loops and coloops

loop



coloop = bridge = isthmus



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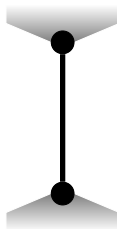
loop



duality



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Contraction and deletion in terms of f

Indicator function of cutset space of G :

$$f : 2^E \rightarrow \{0, 1\}$$

For **contraction** and **deletion** of some $e \in E$:

Indicator functions of cutset spaces of ...

G/e

$$f // e : 2^{E \setminus \{e\}} \rightarrow \{0, 1\}$$

$$f // e(X) = \frac{f(X)}{f(\emptyset)}$$

$G \setminus e$

$$f \parallel e : 2^{E \setminus \{e\}} \rightarrow \{0, 1\}$$

$$f \parallel e(X) = \frac{f(X) + f(X \cup \{e\})}{f(\emptyset) + f(\{e\})}$$

Interpolating between contraction and deletion

(GF, 2004)

For $e \in E$, $X \subseteq E \setminus \{e\}$:

Contraction

$$(f//e)(X)$$

$$\frac{f(X)}{f(\emptyset)}$$

Deletion

$$(f \parallel e)(X)$$

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λ -minor

$$(f \parallel_{\lambda} e)(X)$$

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For $e \in E$, $X \subseteq E \setminus \{e\}$:

Contraction

($\lambda = 0$)

$(f // e)(X)$

$$\frac{f(X)}{f(\emptyset)}$$

λ -minor

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$$\frac{f(X) + \lambda f(X \cup \{e\})}{f(\emptyset) + \lambda f(\{e\})}$$

Deletion

($\lambda = 1$)

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Interpolating between contraction and deletion

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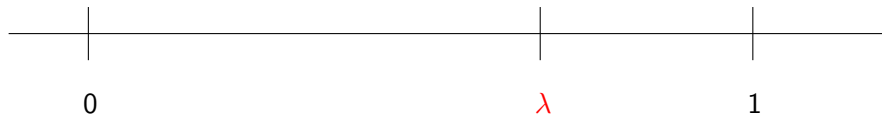
$$\frac{f(X) + \lambda f(X \cup \{e\})}{f(\emptyset) + \lambda f(\{e\})}$$

Deletion

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Duality, contraction and deletion

Duality between contraction and deletion can be extended (GF, 2004).

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$$\widehat{f \parallel_{\lambda} e} = \widehat{f} \parallel_{\lambda^*} e$$

(For binary functions, duality = Hadamard transform (GF, 1993).)

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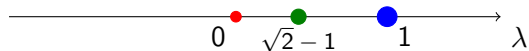
$$\widehat{f \parallel_{\lambda} e} = \widehat{f} \parallel_{\lambda^*} e$$

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Fixed points:

$$\lambda = \pm\sqrt{2} - 1$$

From λ to μ

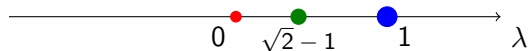


Duality:

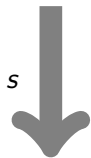
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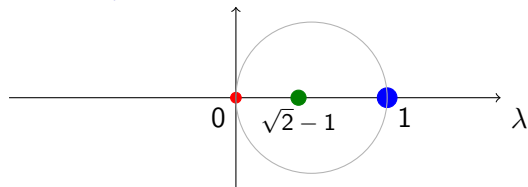


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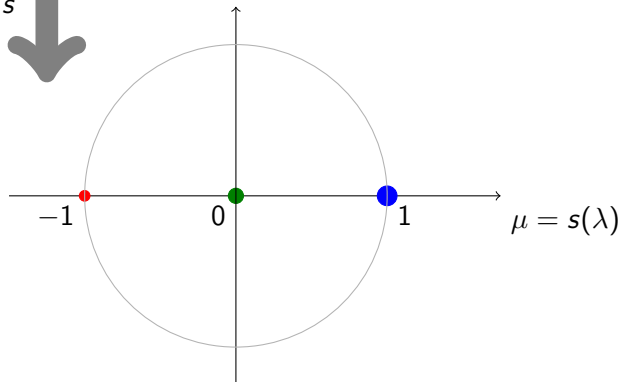
$$\mu^* = -\mu$$

From λ to μ



Duality:

$$\lambda^* = \frac{1 - \lambda}{1 + \lambda}$$



$$\mu^* = -\mu$$

From λ to μ

$$\mu = s(\lambda) := -(3 + 2\sqrt{2}) \frac{\sqrt{2} - 1 - \lambda}{\sqrt{2} + 1 + \lambda}$$

$$\lambda = s^{-1}(\mu) := \frac{1 + \mu}{\sqrt{2} + 1 - (\sqrt{2} - 1)\mu}$$

Notation:

$$f \parallel_{[\mu]} e := f \parallel_{s^{-1}(\mu)} e$$

The transform $L^{[\mu]}$

$$\begin{aligned}(L^{[\mu]}f)(V) &= (2\sqrt{2})^{-|E|} \times \\ &\sum_{X \subseteq E} (\sqrt{2} - 1 + (\sqrt{2} + 1)\mu)^{|X \cap V|} \\ &\quad \cdot (1 - \mu)^{|X \setminus V| + |V \setminus X|} \\ &\quad \cdot (\sqrt{2} + 1 + (\sqrt{2} - 1)\mu)^{|E \setminus (X \cup V)|} f(X)\end{aligned}$$

Matrix representation:

$$M(\mu) = \frac{1}{2\sqrt{2}} \begin{pmatrix} \sqrt{2} + 1 + (\sqrt{2} - 1)\mu & 1 - \mu \\ 1 - \mu & \sqrt{2} - 1 + (\sqrt{2} + 1)\mu \end{pmatrix},$$

$$L^{[\mu]} \mathbf{f} = M(\mu)^{\otimes m} \mathbf{f} \quad (\text{uses } m\text{-th Kronecker power})$$

Special cases:

$\mu = 1$: identity transform

$\mu = -1$: $\sqrt{2}^{|E|} \times$ Hadamard transform (duality)

$\mu = \omega := e^{i2\pi/3}$: some kind of “trianlity”

Properties of the transforms

Composition of transforms \longleftrightarrow multiplication of their parameters:

$$\mathcal{L}^{[\mu_1]} \mathcal{L}^{[\mu_2]} = \mathcal{L}^{[\mu_1 \mu_2]}$$

Also have generalisations of Plancherel's and Parseval's theorems.

$[\mu]$ -minors

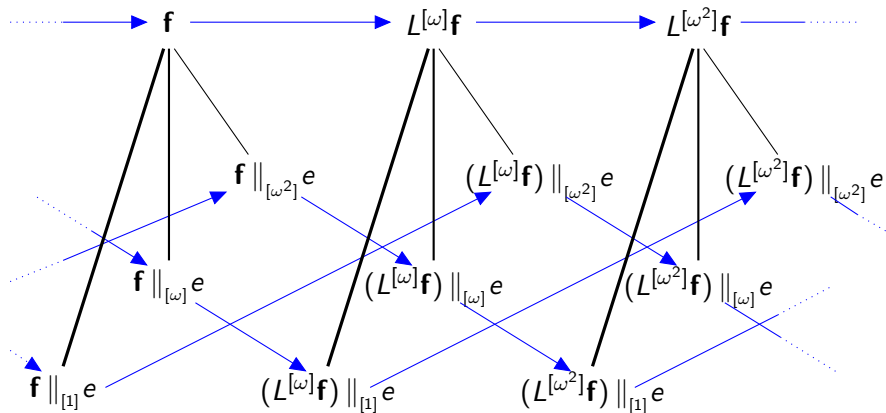
Theorem

$$(L^{[\mu_1]} f) \parallel_{[\mu_2/\mu_1]} e = \text{ScalingFactor}(f, \mu_1, \mu_2) \cdot L^{[\mu_1]}(f \parallel_{[\mu_2]} e)$$

Up to constant factors:

$$\begin{array}{ccc} f & \xrightarrow{L^{[\mu_1]}} & L^{[\mu_1]} f \\ \downarrow [\mu_2]\text{-minor} & & \downarrow [\mu_2/\mu_1]\text{-minor} \\ f \parallel_{[\mu_2]} e & \xrightarrow{L^{[\mu_1]}} & \end{array}$$

$[\omega]$ -minors



Alternating dimaps

Alternating dimap (Tutte, 1948):

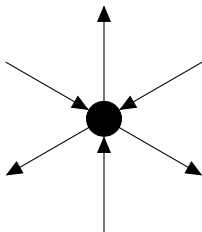
- ▶ directed graph without isolated vertices,
- ▶ 2-cell embedded in a disjoint union of orientable 2-manifolds,
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So vertices look like this:

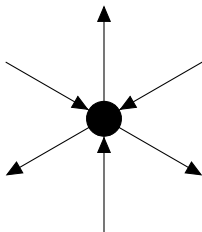


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Genus $\gamma(G)$ of an alternating dimap G :

$$V - E + F = 2(k(G) - \gamma(G))$$

Alternating dimaps

Three special partitions of $E(G)$:

- *clockwise faces*
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Each defines a permutation of $E(G)$. These permutations satisfy

$$\sigma_i \sigma_c \sigma_a = 1$$

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Construction of trial map:

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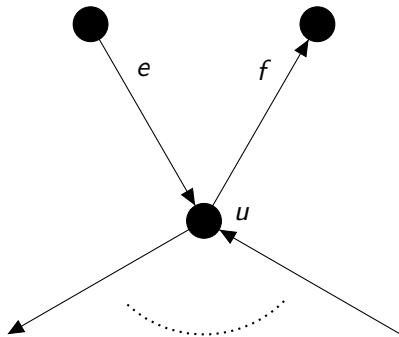
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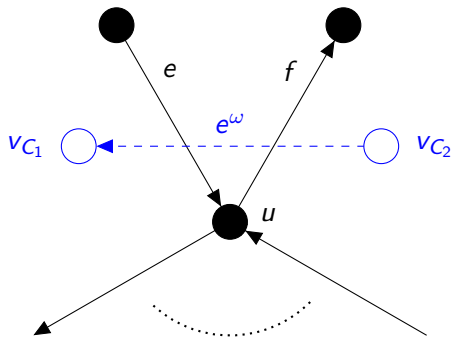


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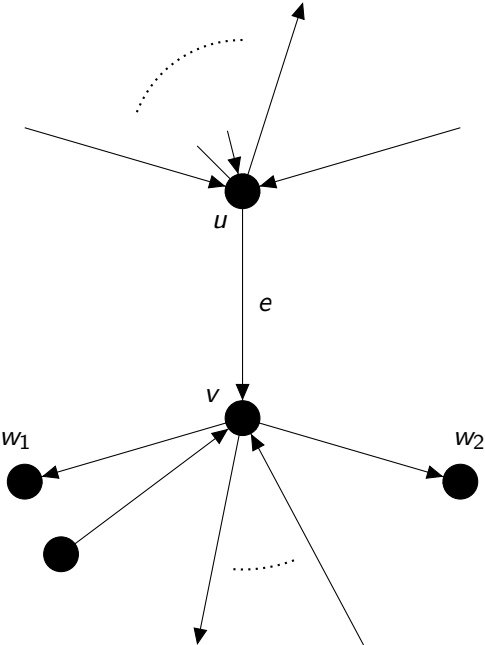
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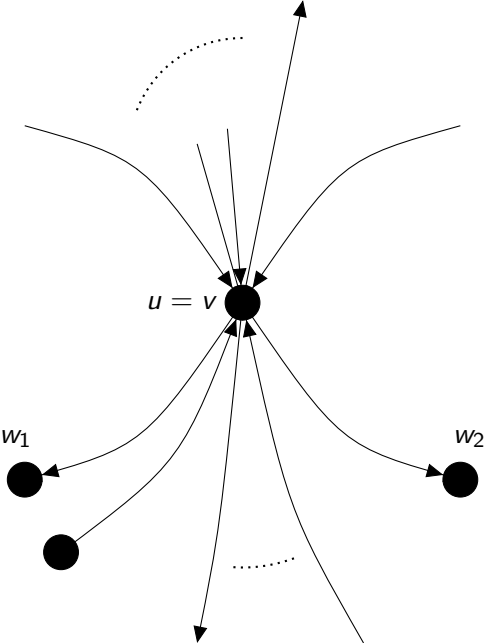
Minor operations

G



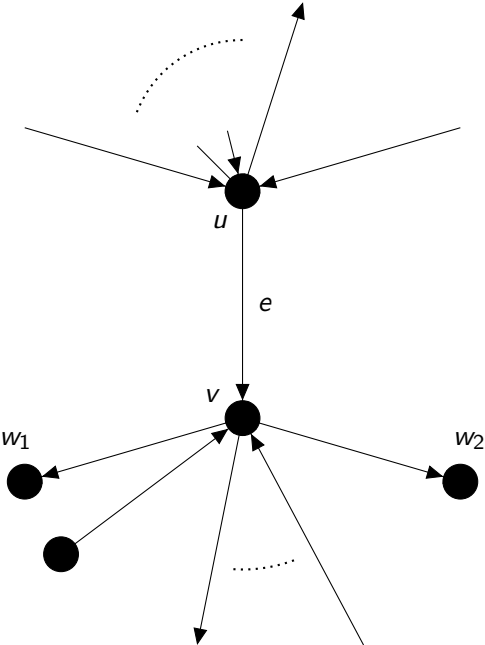
Minor operations

$G[1]e$



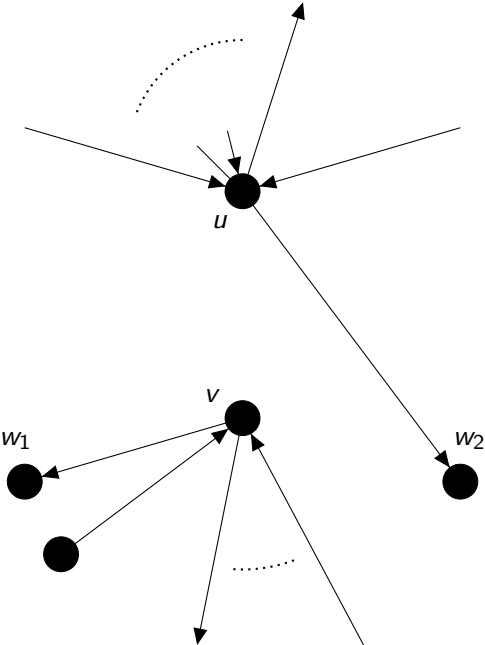
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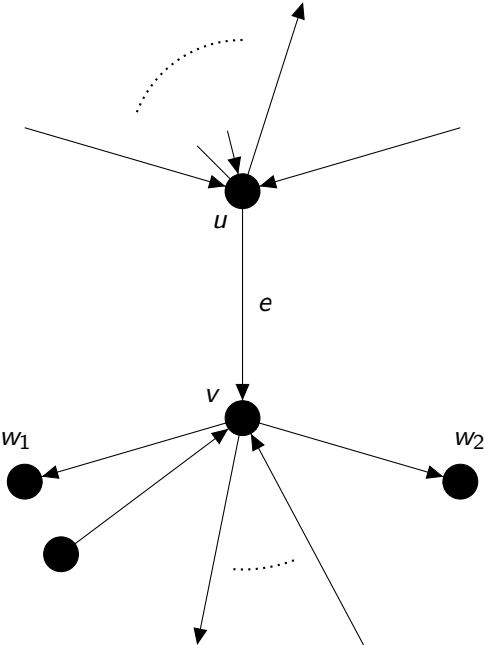
Minor operations

$G[\omega]e$



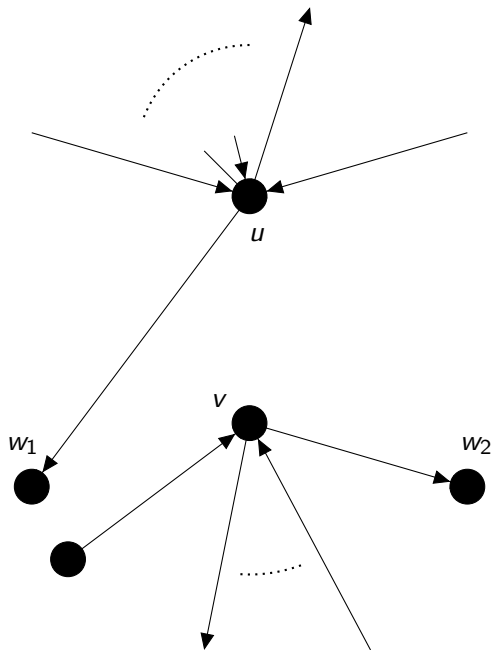
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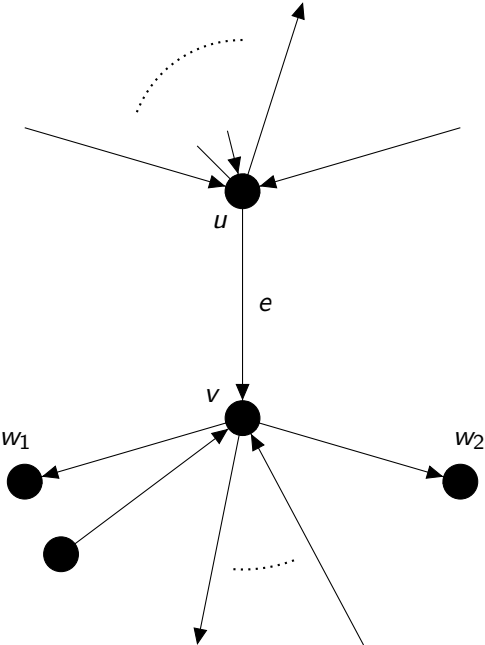
Minor operations

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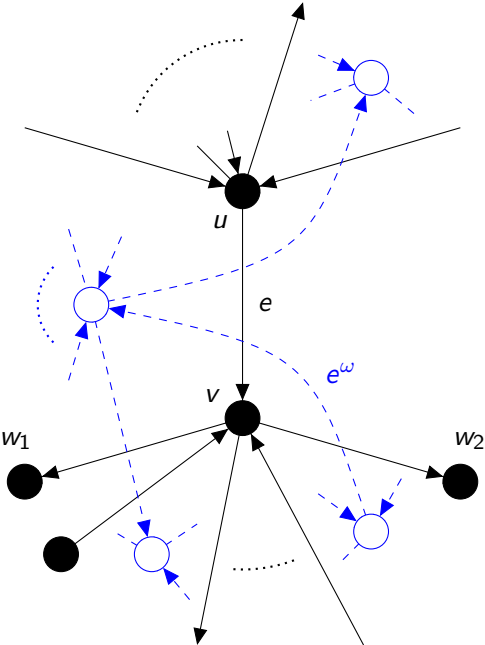
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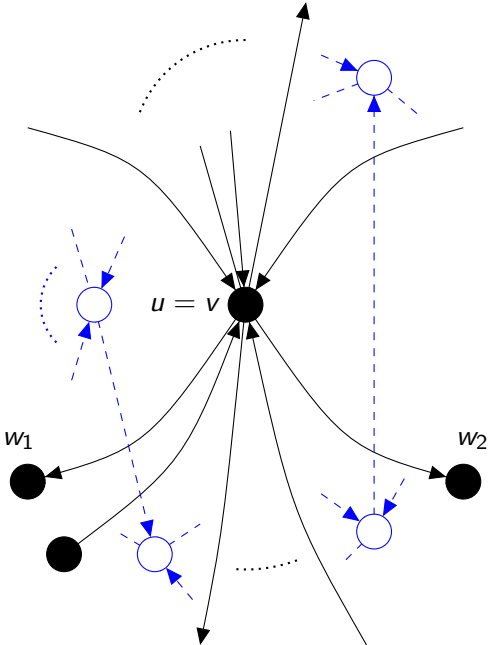
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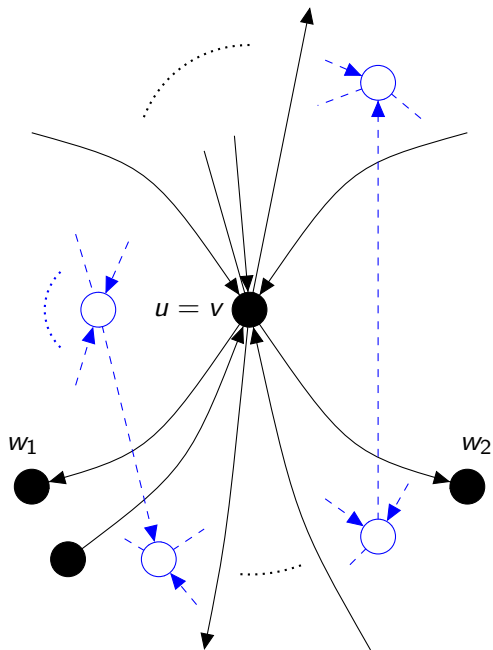
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Minor operations

$$(G[1]e)^\omega = G^\omega[\omega^2]e^\omega$$



Minor operations

$$G^\omega[1]e^\omega = (G[\omega]e)^\omega,$$

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$$G^\omega[\omega^2]e^\omega = (G[1]e)^\omega,$$

$$G^{\omega^2}[1]e^{\omega^2} = (G[\omega^2]e)^{\omega^2},$$

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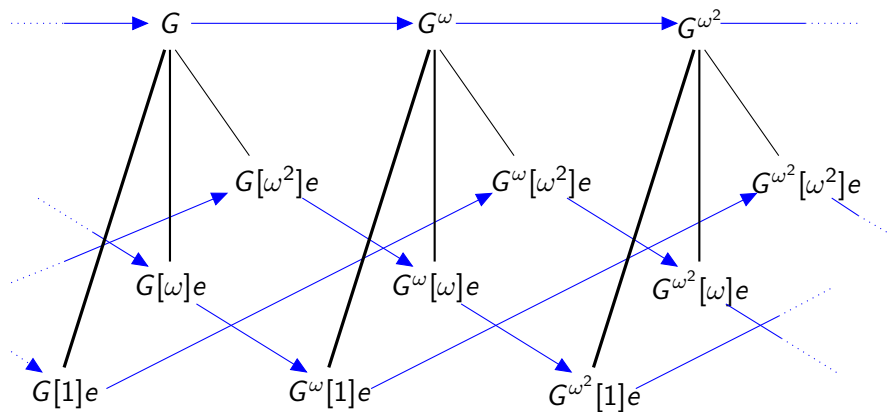
Theorem

If $e \in E(G)$ and $\mu, \nu \in \{1, \omega, \omega^2\}$ then

$$G^\mu[\nu]e^\omega = (G[\mu\nu]e)^\mu.$$

Same pattern as established for generalised minor operations on binary functions (GF, 2008/2013...).

Minor operations



Relationships

triangulated triangle



alternating dimaps



bicubic map (reduction: Tutte 1975)



duality

Eulerian triangulation

Relationships

triangulated triangle



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Eulerian triangulation (reduction, in inverse form ...: Batagelj, 1989)

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(Cavenagh & Lisoněck, 2008)

spherical latin bitrade

Ultraloops, triloops, semiloops

ultraloop



Ultraloops, triloops, semiloops

ultraloop



1-loop

Ultraloops, triloops, semiloops

ω -loop



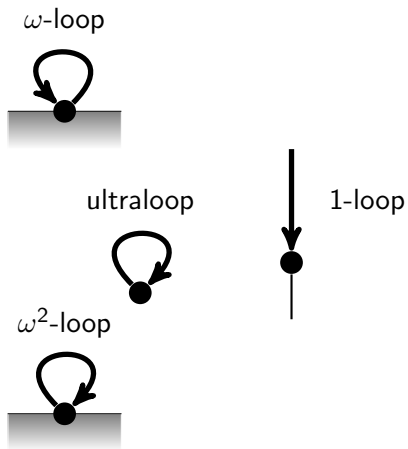
ultraloop



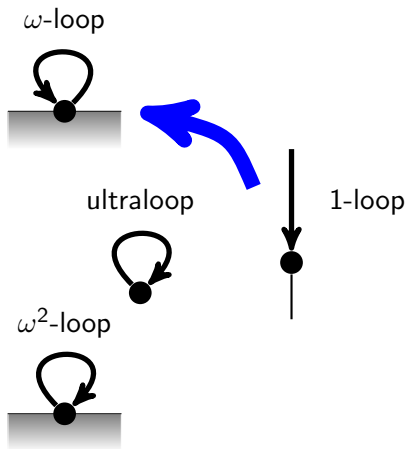
1-loop



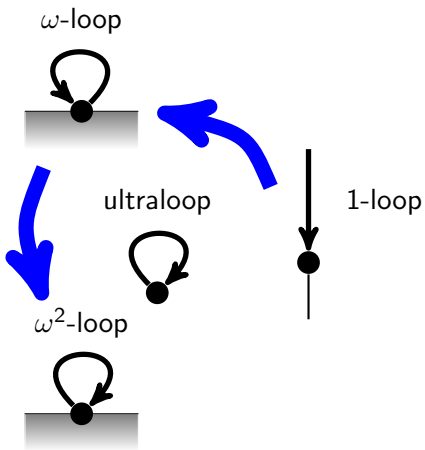
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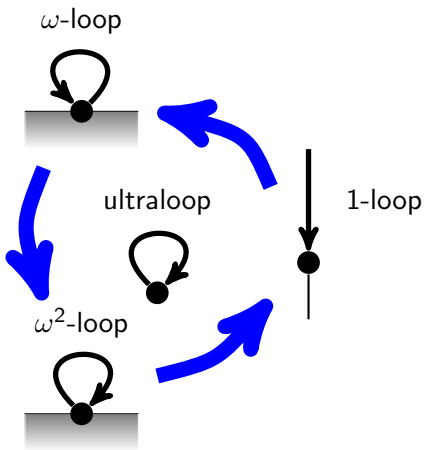
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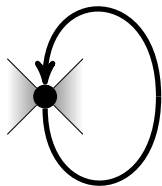


Ultraloops, triloops, semiloops



Ultraloops, triloops, semiloops

1-semiloop



ω -loop



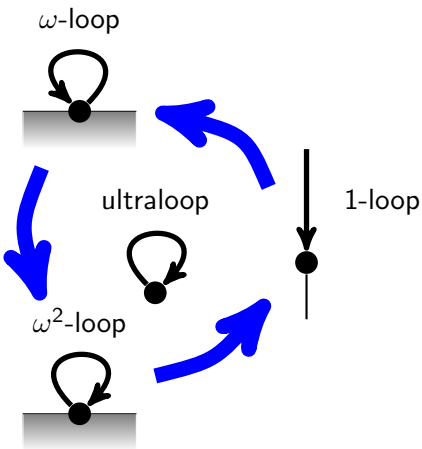
ultraloop



ω^2 -loop

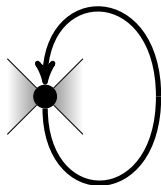


1-loop



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ultraloop



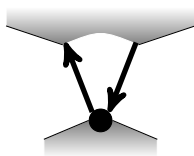
ω^2 -loop



1-loop

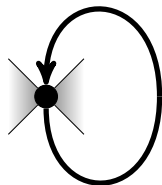


ω -semiloop



Ultraloops, triloops, semiloops

1-semiloop



ω -loop



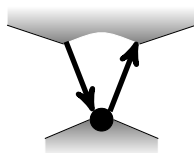
ultraloop



ω^2 -loop



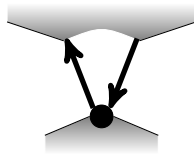
ω^2 -semiloop



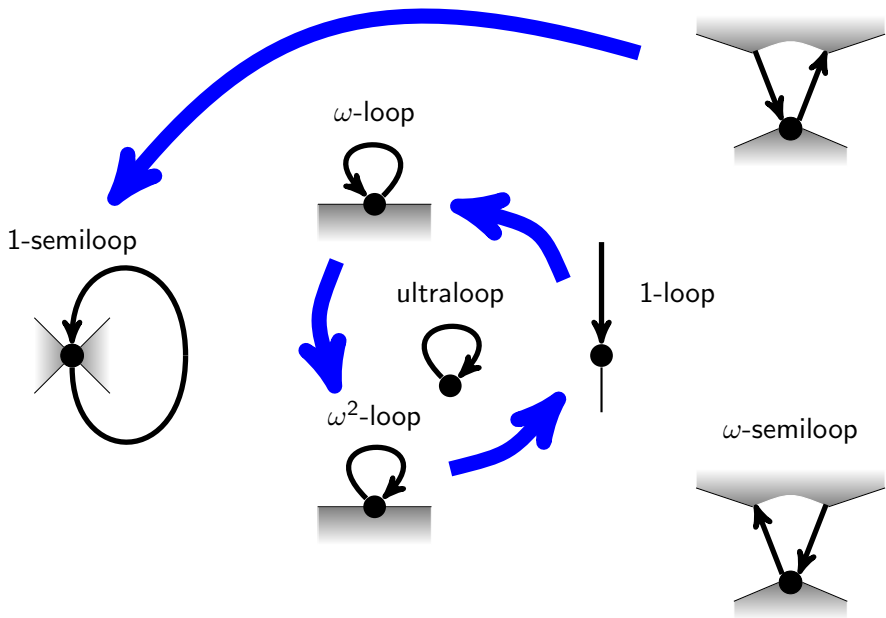
1-loop



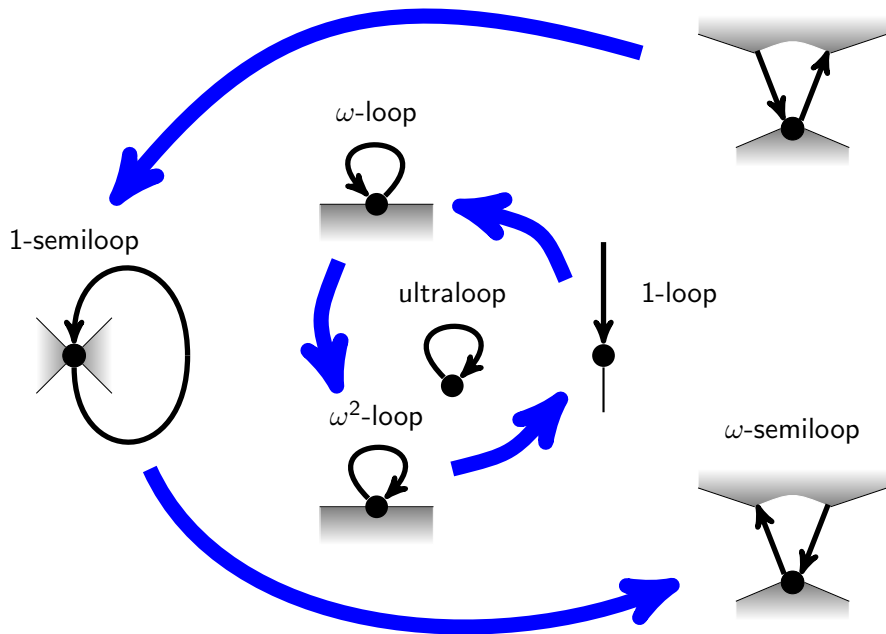
ω -semiloop



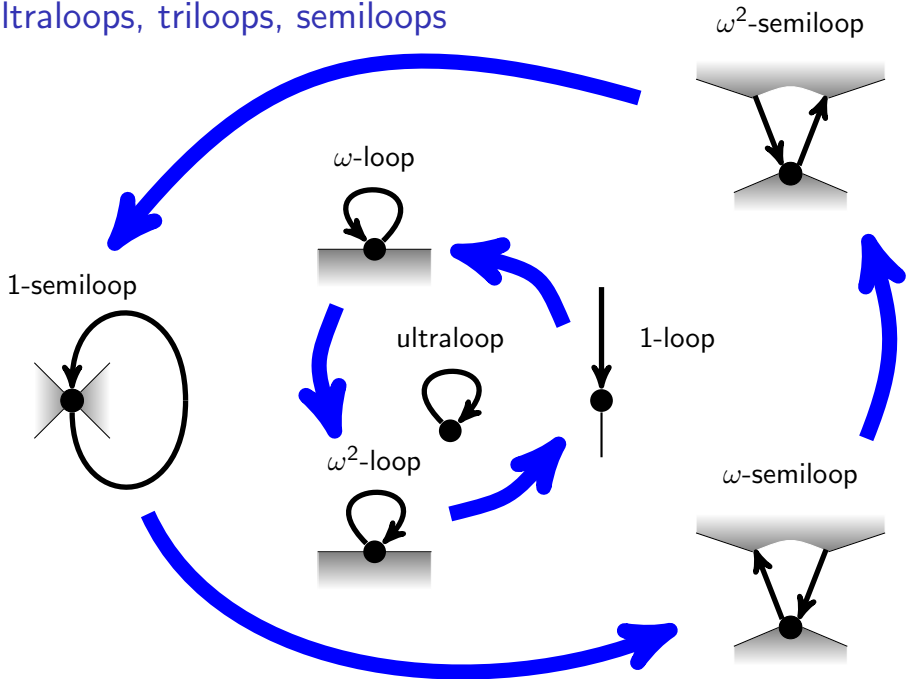
Ultraloops, triloops, semiloops



Ultraloops, triloops, semiloops



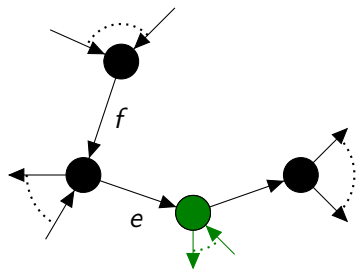
Ultraloops, triloops, semiloops



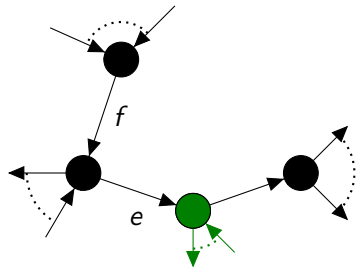
Non-commutativity

Some bad news: sometimes,

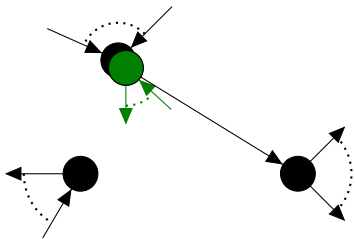
$$G[\mu]e[\nu]f \neq G[\nu]f[\mu]e$$



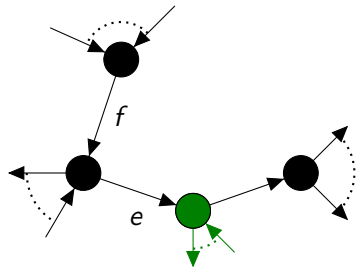
G



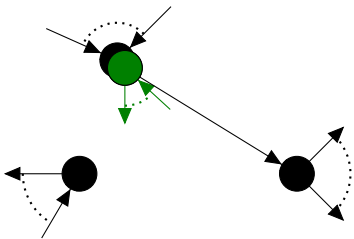
G



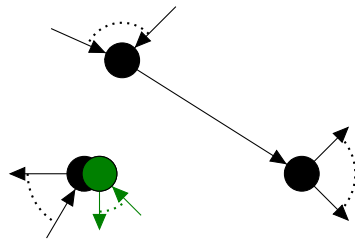
$G[\omega]f[1]e$



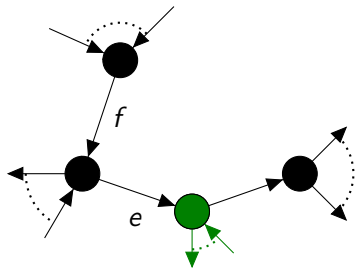
G



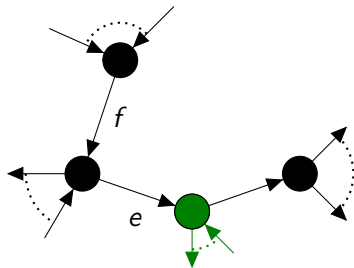
$G[\omega]f[1]e$



$G[1]e[\omega]f$



$$G[\omega]f[1]e \neq G[1]e[\omega]f$$



$$G[\omega]f[1]e \neq G[1]e[\omega]f$$

Theorem

Except for the above situation and its trials, reductions commute.

$$G[\mu]f[\nu]e = G[\nu]e[\mu]f$$

Corollary

If $\mu = \nu$, or one of e, f is a triloop, then reductions commute.

Which alternating dimaps “are” binary functions?

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Not all: for alternating dimaps, reductions do not commute in general, whereas for binary functions, they do.

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Definition

A *strict binary representation* of a minor-closed set \mathcal{A} of alternating dimaps is a triple (F, ε, ν) such that

- (a) $F : \mathcal{A} \rightarrow \{\text{binary functions}\}$
- (b) $\varepsilon = (\varepsilon_G \mid G \in \mathcal{A})$ is a family of bijections
 $\varepsilon_G : E(G) \rightarrow E(F(G))$;
- (c) $\nu \in \mathbb{C}$ with $|\nu| = 1$;
- (d) $F(G^{(\omega)}) \simeq L^{[\omega]}F(G)$ for all $G \in \mathcal{A}$;
- (e) $F(G[\mu]e) \simeq F(G) \parallel_{[\nu\mu]} \varepsilon_G(e)$ for all $G \in \mathcal{A}$, $e \in E(G)$ and $\mu \in \{1, \omega, \omega^2\}$.

Which alternating dimaps are binary functions?

Definitions

C_1 := ultraloop

iC_1 = disjoint union of i ultraloops

$0C_1$ = empty alternating dimap

\mathcal{U}_k = $\{iC_1 \mid i = 0, \dots, k\}$

\mathcal{U}_∞ = $\{iC_1 \mid i \in \mathbb{N} \cup \{0\}\}$

Theorem

If \mathcal{A} is a minor-closed class of alternating dimaps which has a strict binary representation then

- ▶ $\mathcal{A} = \emptyset$, or
- ▶ $\mathcal{A} = \mathcal{U}_k$ for some k , or
- ▶ $\mathcal{A} = \mathcal{U}_\infty$.

Which alternating dimaps are binary functions?

Proof. (Outline) If $\mathcal{A} = \emptyset$: done. So suppose $\mathcal{A} \neq \emptyset$.

Since \mathcal{A} is minor-closed, it must contain the empty alt. dimap $0C_1$. It must be represented by $f : 2^\emptyset \rightarrow \mathcal{C}$ with $f(\emptyset) = 1$, i.e., $\mathbf{f} = (1)$.

If $|\mathcal{A}| = 1$ then we are done. This F gives a strict binary representation, and $\mathcal{A} = \mathcal{U}_0$.

If $|\mathcal{A}| \geq 2$, then it must contain the ultraloop C_1 . Its image $F(C_1)$ is given by

$$F(C_1) = \begin{pmatrix} 1 \\ \sqrt{2} - 1 \end{pmatrix}.$$

Proof: C_1 is self-trial, so $F(C_1)$ must be too. So $F(C_1)$ must be an eigenvector for eigenvalue 1 of the matrix $M(\omega)$.

If $|\mathcal{A}| = 2$ then we are done. This F gives a strict binary representation, and $\mathcal{A} = \{\text{empty, ultraloop}\} = \mathcal{U}_1$.

Which alternating dimaps are binary functions?

Suppose $|\mathcal{A}| \geq 3$. Then \mathcal{A} must have at least one alternating dimap G_2 on two edges.

For any such G_2 , all reductions give the ultraloop C_1 .

So all reductions of $F(G_2)$ give $F(C_1) = \begin{pmatrix} 1 \\ \sqrt{2} - 1 \end{pmatrix}$.

Then show that $F(G_2) = \begin{pmatrix} 1 \\ \sqrt{2} - 1 \end{pmatrix}^{\otimes 2}$.

Therefore $F(G_2)$ is self-trial, so G_2 must be too.

So $G_2 = 2C_1$ (the only self-trial alternating dimap on two edges).

So far, we have at most one alternating dimap in \mathcal{A} with each possible number of edges (0, 1, 2).

Show by induction that \mathcal{A} has at most one member with k edges, and that it is kC_1 , with

$$F(kC_1) = \begin{pmatrix} 1 \\ \sqrt{2} - 1 \end{pmatrix}^{\otimes k}.$$

This is (the guts of) the strict binary representation. □

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