

# Enumeration. Martingales. Random graphs.

Mikhail Isaev

School of Mathematical Sciences, Monash University

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# Introduction

# Expectation of the exponential function

We are interested in estimates for  $\mathbb{E}e^Z$ , where  $Z$  is a complex random variable.

$$\mathbb{E}e^Z \approx e^{\mathbb{E}Z} \quad \text{and} \quad \mathbb{E}e^Z \approx e^{\mathbb{E}Z + \frac{1}{2}\mathbb{E}(Z - \mathbb{E}Z)^2}.$$

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It is clear for the following cases:

- when  $Z$  is small;
- when  $Z = X_1 + \dots + X_n$ , where  $X_1, \dots, X_n$  are independent and small.

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For our purposes we needed:

- when  $Z = f(X_1, \dots, X_n)$ , where  $X_1, \dots, X_n$  are independent;
- when  $Z$  is a complex martingale.

# Random vectors with independent components

Theorem (I., McKay, 2017)

Let  $X = (X_1, \dots, X_n)$  be a random vector with independent components taking values in  $\Omega = \Omega_1 \times \dots \times \Omega_n$ . Let  $f : \Omega \rightarrow \mathbb{C}$ . Suppose, for any  $1 \leq j \neq k \leq n$ ,

$$\sup_{x, x^j} |f(x) - f(x^j)| \leq \alpha,$$

$$\sup_{x, x^j, x^k, x^{jk}} |f(x) - f(x^j) - f(x^k) + f(x^{jk})| \leq \Delta,$$

where the suprema is over all  $x, x^j, x^k, x^{jk} \in \Omega$  such that

- $x, x^j$  and  $x^k, x^{jk}$  differ only in the  $j$ -th coordinate,
- $x, x^k$  and  $x^j, x^{jk}$  differ only in the  $k$ -th coordinate.

(A) If  $\alpha = o(n^{-1/2})$ , then  $\mathbb{E}e^{f(X)} = e^{\mathbb{E}f(X)} (1 + O(n\alpha^2))$ .

(B) If  $\alpha = o(n^{-1/3})$  and  $\Delta = o(n^{-4/3})$ , then

$$\mathbb{E}e^{f(X)} = e^{\mathbb{E}f + \frac{1}{2}\mathbb{E}(f - \mathbb{E}f)^2} \left( 1 + O(n\alpha^3 + n^2\alpha^2\Delta) e^{\frac{1}{2}\text{Var } \mathfrak{S}f(X)} \right).$$

# Random permutations

Theorem (Greenhill, I., McKay, 2017+)

Let  $X = (X_1, \dots, X_n)$  be a uniform random element of  $S_n$  and  $f : S_n \rightarrow \mathbb{C}$ .  
 Suppose, for any distinct  $j, a \in \{1, \dots, n\}$ ,

$$\sup_{\omega \in S_n} |f(\omega) - f(\omega \circ (ja))| \leq \alpha,$$

and, for any distinct  $j, k, a, b \in \{1, \dots, n\}$ ,

$$\sup_{\omega \in S_n} |f(\omega) - f(\omega \circ (ja)) - f(\omega \circ (kb)) + f(\omega \circ (ja)(kb))| \leq \Delta,$$

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# Random subsets

Let  $B_{n,m}$  denote the set of subsets of  $\{1, \dots, n\}$  of size  $m$ .

Theorem (Greenhill, I., McKay, 2017+)

Let  $X$  be a uniform random element of  $B_{n,m}$ ,  $m \leq n/2$ , and  $f : B_{n,m} \rightarrow \mathbb{C}$ . Suppose, for any  $A \in B_{n,m}$  and  $a \in A$ ,  $j \notin A$ ,

$$|f(A) - f(A \oplus \{j, a\})| \leq \alpha,$$

and, for any distinct  $a, b \in A$ ,  $j, k \notin A$ ,

$$|f(A) - f(A \oplus \{j, a\}) - f(A \oplus \{k, b\}) + f(A \oplus \{j, k, a, b\})| \leq \Delta,$$

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- ② Asymptotic normality. Let  $Z = \text{it} \frac{X - \mathbb{E}X}{\sqrt{\text{Var } X}}$ ,  $\phi(t) = \mathbb{E} \exp\left(\text{it} \frac{X - \mathbb{E}X}{\sqrt{\text{Var } X}}\right)$ . Then,

$$\mathbb{E}e^Z = \phi(t) \quad \text{and} \quad e^{\mathbb{E}Z + \frac{1}{2}\mathbb{E}(Z - \mathbb{E}Z)^2} = e^{-t^2/2}.$$

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- ③ Subgraph counts in random graphs with given degrees.
- ④ Asymptotic enumeration by complex-analytic methods.

# Subgraphs counts

## Random variables of our interest

1. Fix a graphical degree sequence  $\vec{d} = (d_1, \dots, d_n)$ .
2. Take a uniform random labelled graph  $\mathcal{G}_{\vec{d}}$  with degrees  $d$ .
3. Count the number  $N_P(\mathcal{G}_{\vec{d}})$  of **occurrences** of a given pattern  $P$ .

We want an **asymptotic** formula for  $\mathbb{E}N_P(\mathcal{G}_{\vec{d}})$  as  $n \rightarrow \infty$ ,  
 where  $P = P(n)$  and  $\vec{d} = \vec{d}(n)$ .

Let  $d$  denote the average degree. Define

$$\lambda = \frac{d}{n-1} \quad \text{and} \quad R = \frac{1}{n} \sum_{j=1}^n (d_j - d)^2.$$



## From probabilities to subgraph counts

We employ formulae from [McKay, 1985], [McKay, 2011] for the probability of a certain pattern  $P$  (subgraph or induced subgraph) to occur at a particular place in our random graph  $\mathcal{G}_{\vec{d}}$ :

$$\Pr(P \text{ occurs in } \mathcal{G}_{\vec{d}} \text{ at location } L) = \text{factor}(P, \lambda) \times e^{\text{stuff}(P, L, \vec{d}) + o(1)}.$$

The only thing to do is to sum up over all possible locations  $L$ .

For example, for copies of a given subgraph it is reduced to estimating

$$\sum_{\sigma \in S_n} e^{f(\sigma)} = n! \mathbb{E} e^{f(X)},$$

where  $X$  is a uniform random permutation.

# Perfect matchings and cycles

Suppose  $\min\{d, n - 1 - d\} \geq cn / \log n$  and  $|d_j - d| \leq n^{1/2+\epsilon}$ . Then, we have

- the expected number of **perfect matchings in  $\mathcal{G}_{\vec{d}}$**  (for even  $n$ ) is

$$\frac{n!}{(n/2)!2^{n/2}} \lambda^{n/2} \exp\left(\frac{1-\lambda}{4\lambda} - \frac{R}{2\lambda^2 n} + o(1)\right);$$

- the expected number of  **$q$ -cycles in  $\mathcal{G}_{\vec{d}}$**  (for  $3 \leq q \leq n$ ) is

$$\frac{n!}{2q(n-q)!} \lambda^q \exp\left(-\frac{(1-\lambda)q(n-q)}{\lambda n^2} + \frac{Rq(n-2q)}{\lambda^2 n^3} + o(1)\right).$$

These expressions for the regular case ( $R = 0$ ) were also given in [McKay, 2011].

# Subgraphs isomorphic to a given one

For a given graph  $H$  with degree sequence  $(h_1, \dots, h_n)$  denote  $\mu_t = \frac{1}{n} \sum_{j=1}^n h_j^t$ .

Theorem 1 (Greenhill, I., McKay, 2017+).

Suppose  $\min\{d, n-1-d\} \geq cn/\log n$  and  $|d_j - d| \leq n^{1/2+\epsilon}$ . Let  $H$  be a graph with degrees  $(h_1, \dots, h_n)$  and  $m = O(n^{1+2\epsilon})$  edges such that, for all  $1 \leq j \leq n$ ,

$$h_j = O(n^{1/2+\epsilon}), \quad \frac{(d_j - d)^3 \mu_3}{\lambda^3} = o(n^2).$$

Then, the expected number of subgraphs in  $\mathcal{G}_{\vec{d}}$  isomorphic to  $H$  is

$$\frac{n!}{|\text{Aut}(H)|} \lambda^m \exp \left( -\frac{1-\lambda}{4\lambda} (2\mu_2 - \mu_1^2 - 2\mu_1) + \frac{R}{2\lambda^2 n} (\mu_2 - \mu_1^2 - \mu_1) \right. \\ \left. - \frac{1-\lambda^2}{6\lambda^2 n} \mu_3 - \frac{1-\lambda}{\lambda n^2} \sum_{j,k \in E(H)} h_j h_k + o(1) \right),$$

where  $\text{Aut}(H)$  is the automorphism group of  $H$ .

A similar result holds for the number of induced copies of  $H$ .

# Spanning trees

Let  $\tau_{\vec{d}}$  denote the number of **spanning trees in  $\mathcal{G}_{\vec{d}}$** .

Theorem (Greenhill, I., McKay, 2017+)

Suppose  $\min\{d, n - 1 - d\} \geq cn / \log n$  and  $|d_j - d| \leq n^{1/2+\varepsilon}$ , then

$$\mathbb{E}\tau_{\vec{d}} = n^{n-2} \lambda^{n-1} \exp\left(-\frac{1-\lambda}{2\lambda} - \frac{R}{2\lambda^2 n} + o(1)\right).$$

Theorem (Greenhill, I., McKay, Kwan, 2017)

Suppose that number of edges is at least  $n + \frac{1}{2}d_{\max}^4$  (i.e.  $d_{\max}^4 \leq (d-2)n$ ), then

$$\mathbb{E}\tau_{\vec{d}} = \frac{(d-1)^{1/2}}{(d-2)^{3/2}n} \left(\prod_{j=1}^n d_j\right) \left(\frac{(d-1)^{d-1}}{d^{d/2}(d-2)^{d/2-1}}\right)^n \times \\ \exp\left(\frac{6d^2 - 14d + 7}{4(d-1)^2} + \frac{R}{2(d-1)^3} + \frac{(2d^2 - 4d + 1)R^2}{4(d-1)^3 d^2} + o(1)\right).$$

# Asymptotic enumeration

# Complex-analytic approach

1. We write combinatorial counts in terms of multivariate generation functions.

Example: the number of  $d$ -regular graphs on  $n$  vertices

$$\text{RG}(n, d) = [x_1^d \cdots x_n^d] \prod_{1 \leq j < k \leq n} (1 + x_j x_k).$$

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$$\text{RG}(n, d) = \frac{1}{(2\pi i)^n} \oint \cdots \oint \frac{\prod_{1 \leq j < k \leq n} (1 + z_j z_k)}{z_1^{d+1} \cdots z_n^{d+1}} dz_1 \cdots, dz_n.$$

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3. By choosing appropriate contours, we approximate the value of the integral which is mostly given by **small neighbourhoods** of concentration points.



# Integrals we work with

Typically the problem is reduced to the estimation of integrals of the following form:

$$\int_B \mathbf{exp}(-x^T Ax + f(x)) dx = c \mathbb{E} e^{f(X_B)},$$

where  $x = (x_1, \dots, x_n) \in \mathbb{R}^n$ ,  $f$  is a multi-variable polynomial of low degree with complex coefficients and  $X_B$  is a gaussian random variable truncated to  $B$ .

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$$\int_{\mathbf{B}} \exp(-\mathbf{x}^T \mathbf{A} \mathbf{x} + f(\mathbf{x})) d\mathbf{x} = c \mathbb{E} e^{f(\mathbf{X}_{\mathbf{B}})},$$

where  $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{R}^n$ ,  $f$  is a multi-variable polynomial of low degree with complex coefficients and  $\mathbf{X}_{\mathbf{B}}$  is a gaussian random variable truncated to  $\mathbf{B}$ .

Example: if  $nd$  is even and  $\min\{d, n-1-d\}$  grows sufficiently fast as  $n \rightarrow \infty$

$$\text{RG}(n, d) \approx \frac{2(2\pi)^{-n}}{(\lambda^\lambda(1-\lambda)^{1-\lambda}) \binom{n}{2}} \int_{\|\mathbf{x}\|_\infty \leq \frac{n^\epsilon}{(\lambda(1-\lambda)_n)^{1/2}}} \exp\left(\sum_{\ell=2}^{\ell_{\max}} \sum_{j < k} c_\ell (x_j + x_k)^\ell\right) d\mathbf{x},$$

where  $\lambda = d/(n-1)$  is the density of such a graph and  $c_2 = -\frac{1}{2}\lambda(1-\lambda)$ .

# The $\beta$ -model of random graph

A random graph model with independent adjacencies

$$p_{jk} = \frac{e^{\beta_j + \beta_k}}{1 + e^{\beta_j + \beta_k}}.$$

is known as  $\beta$ -model, where  $\beta = (\beta_1, \dots, \beta_n) \in \mathbb{R}^n$ .

Let  $\mathcal{G}_\beta$  denote the  $\beta$ -model where  $\beta$  is the solution of the system

$$\sum_{k:j \neq k} p_{jk} = d_j, \quad (1 \leq j \leq n).$$

for a given degree sequence  $\vec{d} = (d_1, \dots, d_n)$ .

The model  $\mathcal{G}_\beta$  behaves similar to  $\mathcal{G}_{\vec{d}}$  in many ways.

## Some corollaries

Using the formula  $\mathbb{E}e^{f(X)} \approx e^{\mathbb{E}f(X) + \frac{1}{2}\mathbb{E}(f(X) - \mathbb{E}f(X))^2}$  it was obtained in [I., McKay, 2017] that (for the dense case):

- Models  $\mathcal{G}_\beta$  and  $\mathcal{G}_{\vec{d}}$  agree for small subgraph probabilities.

$$P_{\vec{d}}(H^+, H^-) \approx \prod_{jk \in H^+} p_{jk} \prod_{jk \in H^-} (1 - p_{jk}).$$

- Let

$$X_{\vec{d}} = |\mathcal{G}_{\vec{d}} \cap Y| \quad \text{and} \quad X_\beta = |\mathcal{G}_\beta \cap Y|,$$

where  $Y$  is a set of edges. Then, we have

$$\Pr(|X_{\vec{d}} - \mathbb{E}X_\beta| > t|Y|^{1/2}) \leq ce^{-2t} \min\{t, n^{1/6}(\log n)^{-3}\}.$$

# Cumulant expansion

## Moments and cumulants

One could use series  $\mathbb{E}e^{tZ} = \sum_{s=0}^{\infty} \frac{t^s}{s!} \mathbb{E}Z^s$ . However, for  $Z = f(X_B)$ , we would need to estimate **a big number of moments, before they become negligible.**

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Instead, we employ the cumulant expansion

$$\mathbb{E}e^{tZ} = \exp\left(\sum_{r=1}^{\infty} \frac{t^r}{r!} \kappa_r(Z)\right).$$

Recall that cumulants  $\kappa_r(Z)$  could be defined by

$$\kappa_r(Z) = \sum_{\pi} (|\pi| - 1)! (-1)^{|\pi|-1} \prod_{P \in \pi} \mathbb{E}Z^{|P|},$$

where the sum is over all partitions  $\pi$  of  $\{1, \dots, n\}$ . For example,

$$\kappa_1(Z) = \mathbb{E}Z, \quad \kappa_2(Z) = \mathbb{E}Z^2 - (\mathbb{E}Z)^2 = \mathbb{E}(Z - \mathbb{E}Z)^2,$$

$$\kappa_3(Z) = \mathbb{E}Z^3 - 3(\mathbb{E}Z^2)(\mathbb{E}Z) + 2(\mathbb{E}Z)^3 = \mathbb{E}(Z - \mathbb{E}Z)^3,$$

$$\kappa_4(Z) = \mathbb{E}Z^4 - 4(\mathbb{E}Z^3)(\mathbb{E}Z) - 3(\mathbb{E}Z^2)^2 + 12(\mathbb{E}Z^2)(\mathbb{E}Z)^2 - 6(\mathbb{E}Z)^4.$$

## Overview of specifics

For enumeration of regular graphs we have  $f(x) = \sum_{\ell=3}^{\ell_{\max}} \sum_{j < k} c_{\ell}(x_j + x_k)^{\ell}$ .

Main difficulties:

- $Z = f(X_B)$  is not a sum of independent random variables.
- The dimension  $n$  is the parameter that goes to infinity.
- The multi-variable polynomial  $f$  is complex-valued.



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In particular, the following bound of error terms is not good enough:

$$|\mathbb{E}(We^Z)| \leq \|W\|_{\infty} \mathbb{E}|e^Z| = \|W\|_{\infty} \mathbb{E}e^{\Re Z},$$

because  $\mathbb{E}e^{\Re Z}$  could be much bigger than  $\mathbb{E}e^Z$ .

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In all previous works the method is limited to the “dense” range (where  $\mathbb{E}e^{\Re Z}$  is of the same order as  $\mathbb{E}e^Z$ ) and the formulae  $\mathbb{E}e^Z \approx e^{\mathbb{E}Z}$  and  $\mathbb{E}e^Z \approx e^{\mathbb{E}Z + \frac{1}{2}\text{V}Z}$  are used.

# The gap between sparse and dense

Conjecture [McKay, Wormald, 1991]

Suppose  $dn$  is even and  $0 < d < n - 1$ , then the number of  $d$ -regular graphs on  $n$  vertices is

$$\text{RG}(n, d) = \left( \lambda^\lambda (1 - \lambda)^{1-\lambda} \right)^{\binom{n}{2}} \binom{n-1}{d}^n \sqrt{2} e^{1/4+o(1)}$$

as  $n \rightarrow \infty$ , where  $\lambda = \frac{d}{n-1}$ .

We know it is true for the following cases:

- $\min\{d, n-1-d\} = o(n^{1/2})$  (McKay and Wormald, 1991).
- $\min\{d, n-1-d\} \geq c \frac{n}{\log n}$  (McKay and Wormald, 1990).

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# Our results for $d$ -regular graphs

Let  $L = \lambda(1 - \lambda) = \frac{d(n-1-d)}{(n-1)^2}$  and define  $\delta(n, d)$  by

$$\text{RG}(n, d) = \left( \lambda^\lambda (1 - \lambda)^{1-\lambda} \right)^{\binom{n}{2}} \binom{n-1}{d}^n \sqrt{2} e^{1/4 + \delta(n, d)}.$$

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Theorem (I., McKay)

Suppose  $dn$  is even and  $n^{1/7+\epsilon} \leq d \leq \frac{n-1}{2}$ , then (up to the error term  $O(n/d^7)$ )

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$$\begin{aligned} \delta(n, d) &= O(n^{-1}) = -\frac{1}{4n} + O\left(\frac{1}{dn}\right) \\ &= -\frac{1}{4n} + \frac{2 - 23L}{24Ln^2} + O\left(\frac{1}{dn^2}\right) \end{aligned}$$



# Our results for $d$ -regular graphs

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$$\text{RG}(n, d) = \left(\lambda^\lambda(1 - \lambda)^{1-\lambda}\right)^{\binom{n}{2}} \binom{n-1}{d}^n \sqrt{2} e^{1/4 + \delta(n, d)}.$$

Theorem (I., McKay)

Suppose  $dn$  is even and  $n^{1/7+\epsilon} \leq d \leq \frac{n-1}{2}$ , then (up to the error term  $O(n/d^7)$ )

$$\begin{aligned} \delta(n, d) &= O(n^{-1}) = -\frac{1}{4n} + O\left(\frac{1}{dn}\right) \\ &= -\frac{1}{4n} + \frac{2 - 23L}{24Ln^2} + O\left(\frac{1}{dn^2}\right) \\ &= -\frac{1}{4n} + \frac{2 - 23L}{24Ln^2} + \frac{22 - 129L}{24Ln^3} + O\left(\frac{1}{d^2n^2}\right) \end{aligned}$$

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# Regular graphs on 4, 5, 6 vertices

$$\text{RG}(4,1) = \text{RG}(4,2) = 3.$$

$$\text{RG}(5,2) = 12.$$

$$\text{RG}(6,1) = \text{RG}(6,4) = 15.$$

$$\text{RG}(6,2) = \text{RG}(6,3) = 70.$$

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$$\text{RG}(4,1) = \text{RG}(4,2) = 3.$$

$$F_0 = 3.23, F_1 = 3.03, F_2 = 2.92, F_3 = 2.87, F_4 = 2.84.$$

$$\text{RG}(5,2) = 12.$$

$$F_0 = 13.79, F_1 = 13.11, F_2 = 12.79, F_3 = 12.61, F_4 = 12.50.$$

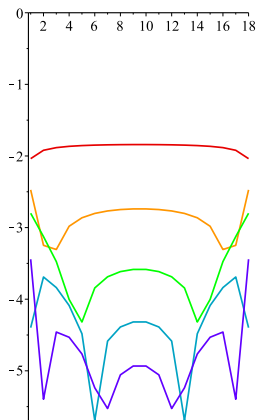
$$\text{RG}(6,1) = \text{RG}(6,4) = 15.$$

$$F_0 = 15.60, F_1 = 14.96, F_2 = 14.78, F_3 = 14.80, F_4 = 14.91.$$

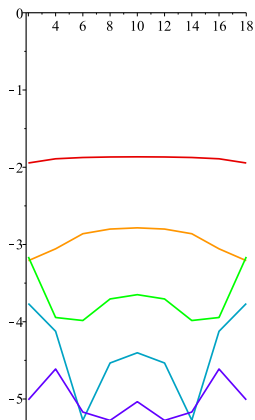
$$\text{RG}(6,2) = \text{RG}(6,3) = 70.$$

$$F_0 = 74.96, F_1 = 71.90, F_2 = 70.68, F_3 = 70.18, F_4 = 69.92.$$

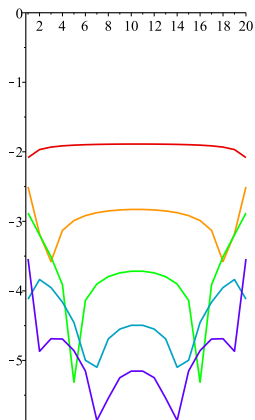
# Plots of $\log_{10} \left( \frac{|F_x - RG|}{RG} \right)$ for $n = 20, 21, 22$



$n = 20.$

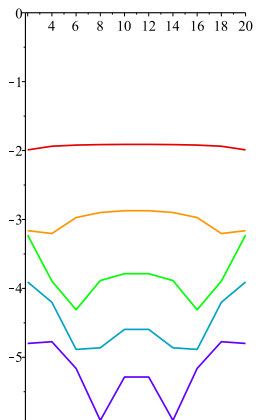
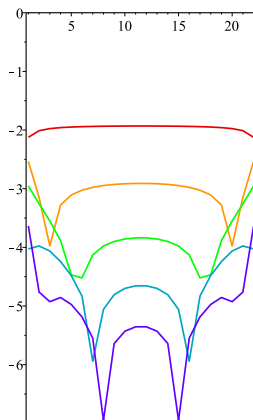
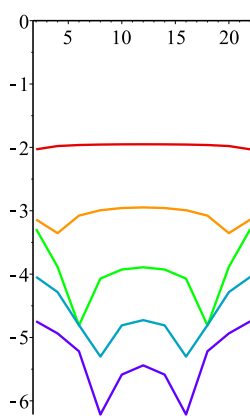


$n = 21.$



$n = 22.$

# Plots of $\log_{10} \left( \frac{|F_x - RG|}{RG} \right)$ for $n = 23, 24, 25$

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# Perfect matchings

Note that  $\text{RG}(2k, 1) = (2k - 1)(2k - 3) \cdots 1 = \frac{(2k)!}{2^k k!}$ .

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$$F_0 = 2.73002 \cdot 10^{78}.$$

$$F_1 = 2.72320 \cdot 10^{78}.$$

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**Conclusion:** the cumulant expansion not only helps to extend the range of complex analytic-methods, but also give more accurate approximations.

Thank you for your attention!