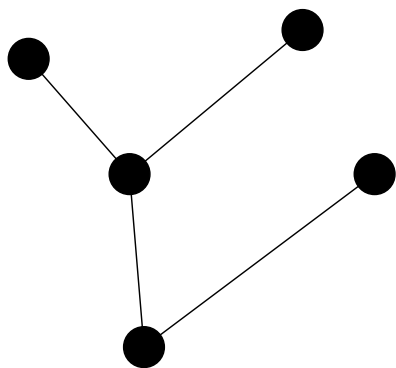


# **Euler Characteristic**

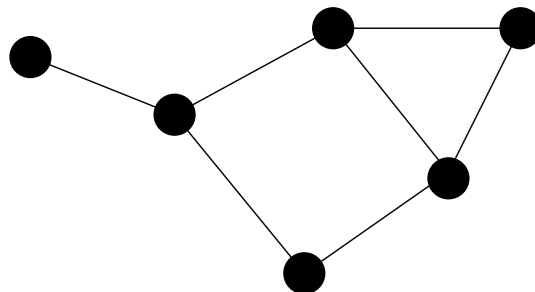
**Rebecca Robinson**

**May 15, 2007**

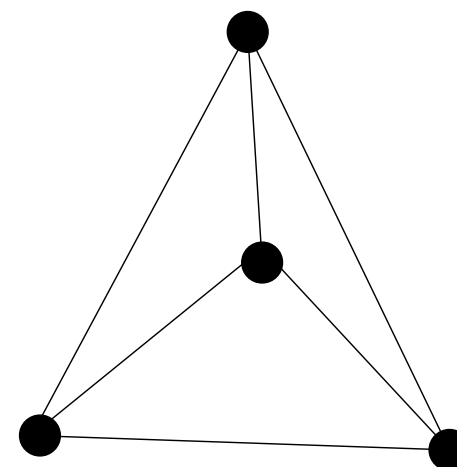
# 1 Planar graphs



$$v = 5, e = 4, f = 1$$
$$v - e + f = 2$$



$$v = 6, e = 7, f = 3$$
$$v - e + f = 2$$



$$v = 4, e = 6, f = 4$$
$$v - e + f = 2$$

*Euler characteristic:*  $\chi = v - e + f$

If a finite, connected, planar graph is drawn in the plane without any edge intersections, and:

- $v$  is the number of vertices,
- $e$  is the number of edges, and
- $f$  is the number of faces

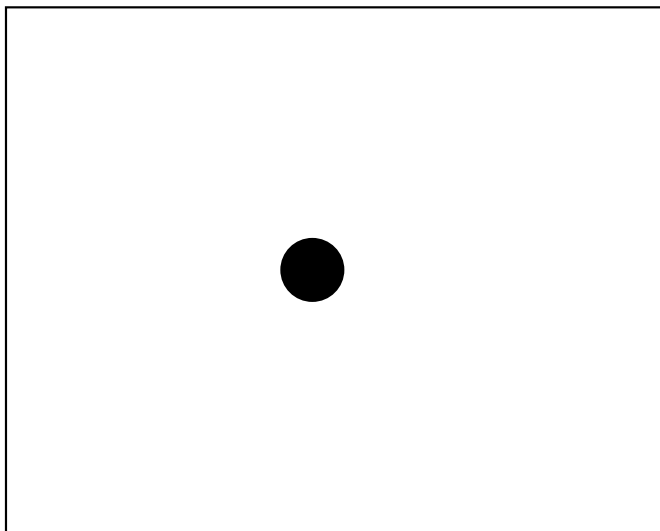
then:

$$\chi = v - e + f = 2$$

ie. the Euler characteristic is 2 for planar surfaces.

Proof.

Start with smallest possible graph:



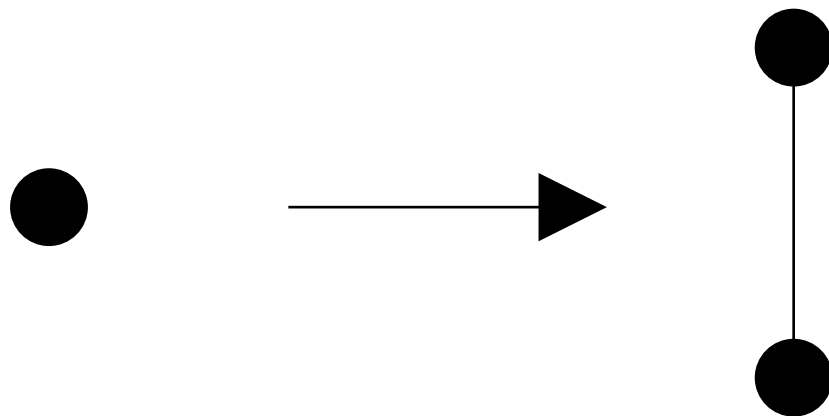
$$v = 1, e = 0, f = 1$$

$$v - e + f = 2$$

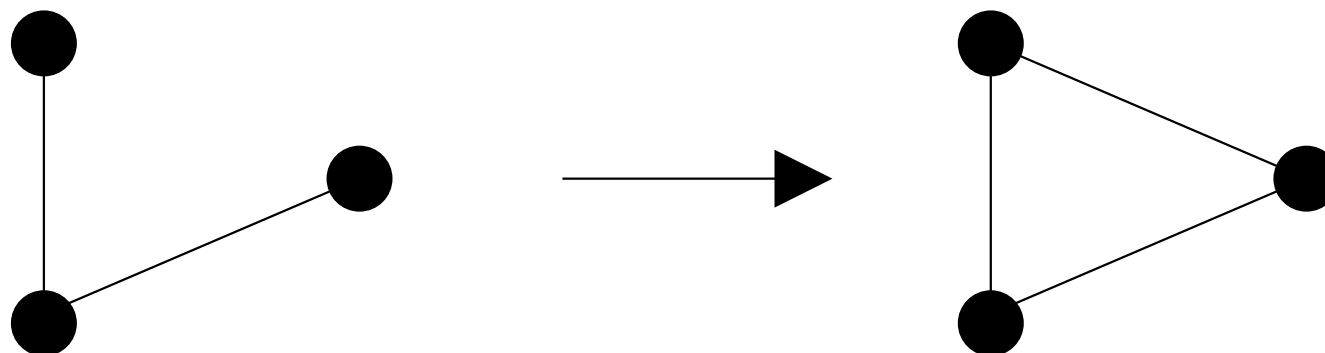
Holds for base case

Increase size of graph:

- either add a new edge and a new vertex, keeping the number of faces the same:



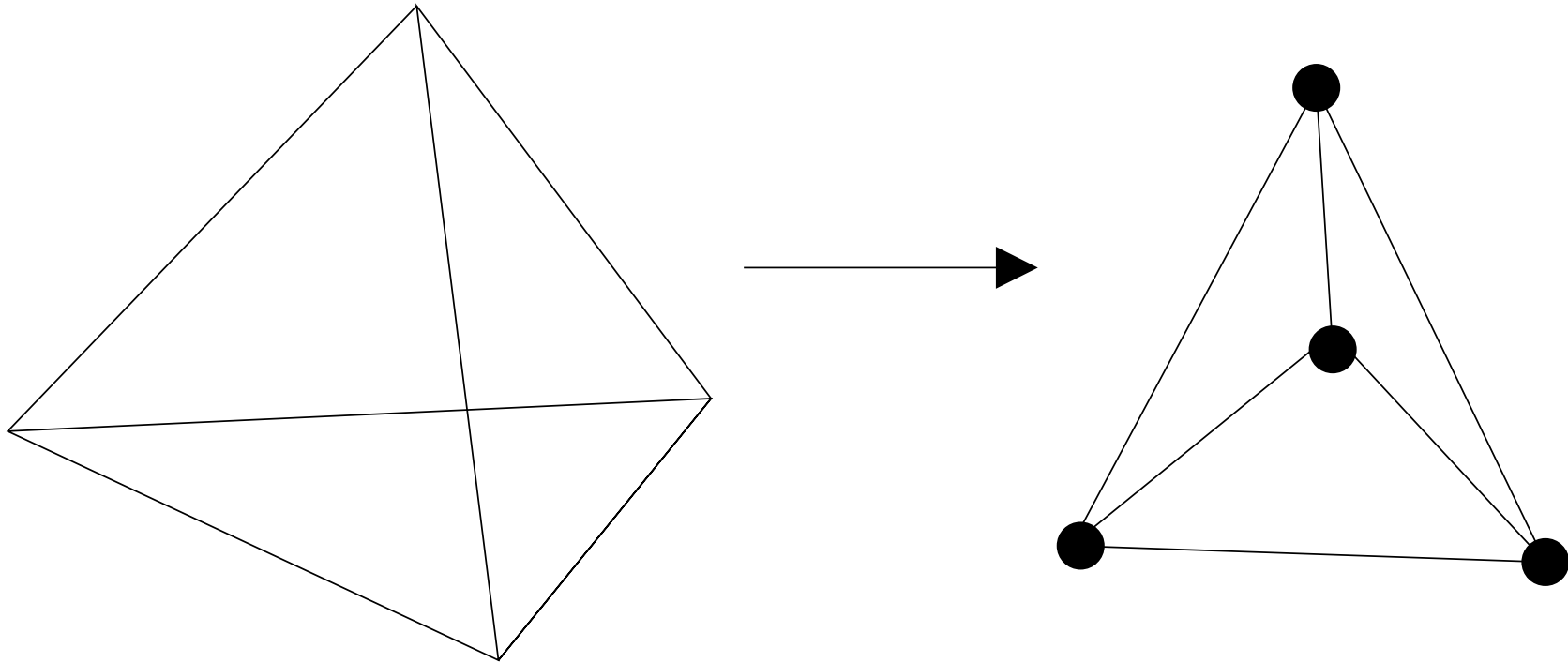
- or add a new edge but no new vertex, thus completing a new cycle and increasing the number of faces:



## 2 Polyhedra

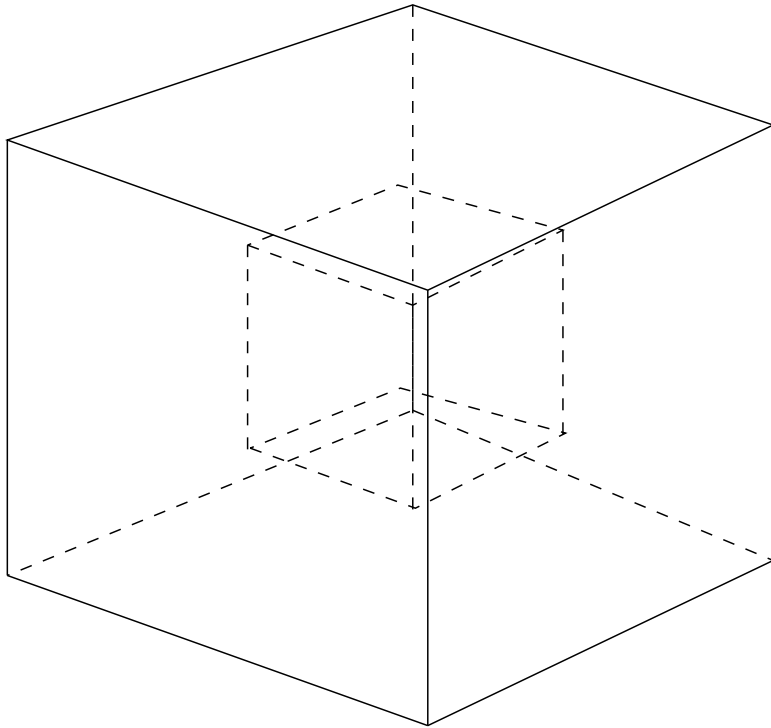
- Euler first noticed this property applied to polyhedra
- He first mentions the formula  $v - e + f = 2$  in a letter to Goldbach in 1750
- Proved the result for convex polyhedra in 1752

- Holds for polyhedra where the vertices, edges and faces correspond to the vertices, edges and faces of a connected, planar graph

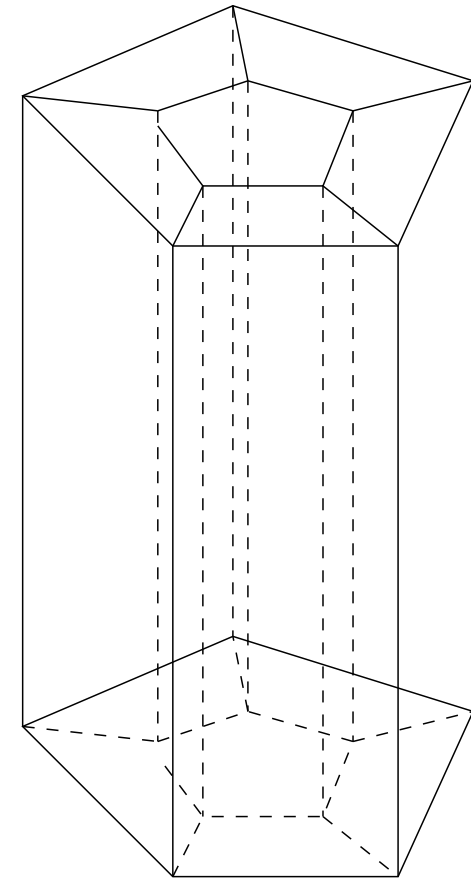




- In 1813 Lhuilier drew attention to polyhedra which did not fit this formula



$$v = 16, e = 24, f = 12$$
$$v - e + f = 4$$



$$v = 20, e = 40, f = 20$$
$$v - e + f = 0$$

**Euler's theorem.** (Von Staudt, 1847) Let  $P$  be a polyhedron which satisfies:

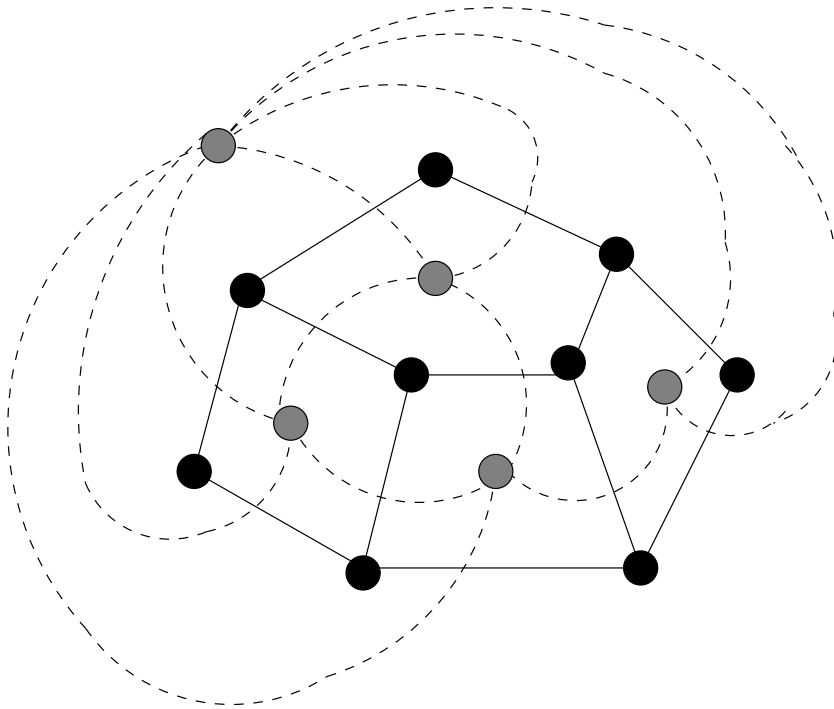
- (a) Any two vertices of  $P$  can be connected by a chain of edges.
- (b) Any loop on  $P$  which is made up of straight line segments (not necessarily edges) separates  $P$  into two pieces.

Then  $v - e + f = 2$  for  $P$ .

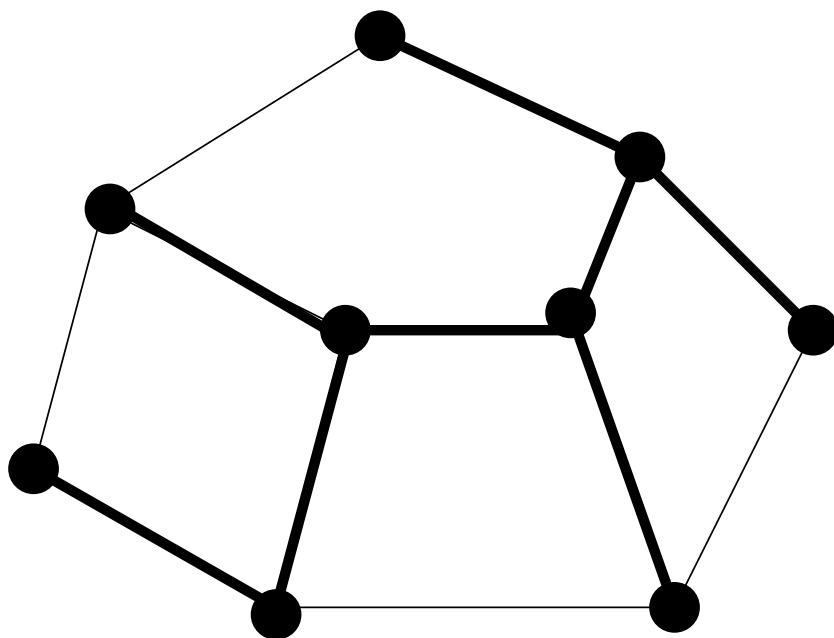
Von Staudt's proof:

For a connected, planar graph  $G$ , define the *dual graph*  $G'$  as follows:

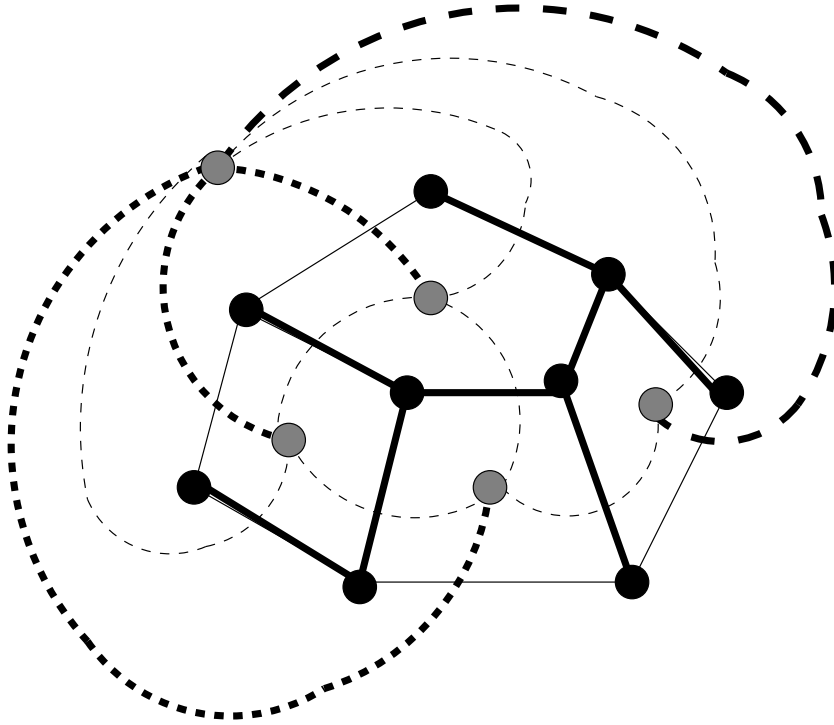
- add a vertex for each face of  $G$ ; and
- add an edge for each edge in  $G$  that separates two neighbouring faces.



Choose a spanning tree  $T$  in  $G$ .



Now look at the edges in the dual graph  $G'$  of  $T'$ 's complement ( $G - T$ ).



The resulting graph  $T'$  is a spanning tree of  $G'$ .

- Number of vertices in any tree = number of edges + 1.

$$|V(T)| - |E(T)| = 1 \text{ and } |V(T')| - |E(T')| = 1$$

$$|V(T)| - [|E(T)| + |E(T')|] + |V(T')| = 2$$

$$|V(T)| = |V(G)|, \text{ since } T \text{ is a spanning tree of } G$$

$$|V(T')| = |F(G)|, \text{ since } T' \text{ is a spanning tree of } G\text{'s dual}$$

$$|E(T)| + |E(T')| = |E(G)|$$

- Therefore  $V - E + F = 2$ .

- *Platonic solid*: a convex, regular polyhedron, i.e. one whose faces are identical and which has the same number of faces around each vertex.
- Euler characteristic can be used to show there are exactly five Platonic solids.

Proof.

Let  $n$  be the number of edges and vertices on each face. Let  $d$  be the degree of each vertex.

$$nF = 2E = dV$$

Rearrange:

$$e = dV/2, f = dV/n$$

By Euler's formula:

$$V - dV/2 + dV/n = 2$$

$$V(2n + 2d - nd) = 4n$$

Since  $n$  and  $V$  are positive:

$$2n + 2d - nd > 0$$

$$(n - 2)(d - 2) < 4$$

Thus there are five possibilities for  $(d, n)$ :

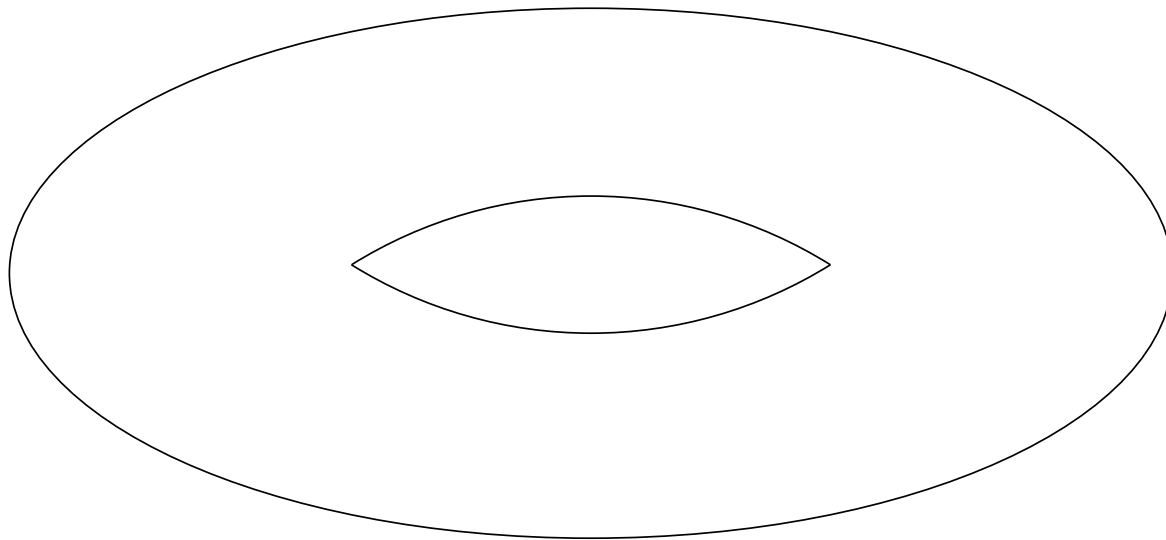
$(3, 3)$  (tetrahedron),  $(3, 4)$  (cube),  $(3, 5)$  (dodecahedron),  $(4, 3)$  (octahedron),  
 $(5, 3)$  (icosahedron).



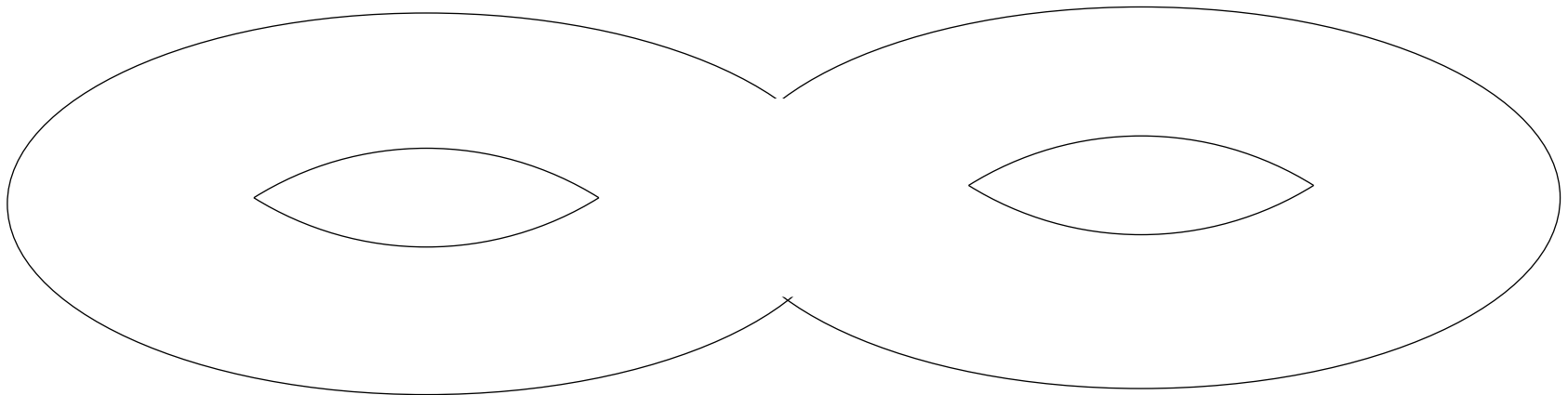
### 3 Non-planar surfaces

- $\chi = v - e + f = 2$  applies for graphs drawn on the plane - what about other surfaces?
- *Genus* of a graph: a number representing the maximum number of cuttings that can be made along a surface without disconnecting it - the number of *handles* of the surface.
- In general:  $\chi = 2 - 2g$ , where  $g$  is the genus of the surface
- Plane has genus 0, so  $2 - 2g = 2$

Torus (genus 1):  $v - e + f = 0$



Double torus (genus 2):  $v - e + f = -2$



- *Topological equivalence*: two surfaces are topologically equivalent (or *homeomorphic*) if one can be ‘deformed’ into the other without cutting or gluing.
- Examples: the sphere is topologically equivalent to any convex polyhedron; a torus is topologically equivalent to a ‘coffee cup’ shape.
- Topologically equivalent surfaces have the same Euler number: the Euler characteristic is called a *topological invariant*